

The number of pancyclic arcs in a k -strong tournament

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A tournament is a digraph, where there is precisely one arc between every pair of distinct vertices. An arc is pancyclic in a digraph D , if it belongs to a cycle of length l , for all $3 \leq l \leq |V(D)|$. Let $p(D)$ denote the number of pancyclic arcs in a digraph D and let $h(D)$ denote the maximum number of pancyclic arcs belonging to the same Hamilton cycle of D . Note that $p(D) \geq h(D)$. Moon showed that $h(T) \geq 3$ for all strong non-trivial tournaments, T , and Havet showed that $h(T) \geq 5$ for all 2-strong tournaments T . We will in this talk show that if T is a k -strong tournament, with $k \geq 2$, then $p(T) \geq nk/2$ and $h(T) \geq (k + 5)/2$. This solves a conjecture by Havet, stating that there exists a constant a_k , such that $p(T) \geq a_k * n$, for all k -strong tournaments, T , with $k \geq 2$. Furthermore the second result gives support for the conjecture $h(T) \geq 2k + 1$, which was also stated by Havet. The previously best known bounds when $k \geq 2$ were $p(T) \geq 2k + 3$ and $h(T) \geq 5$. Furthermore some of the lemma's used in the above proofs immediately imply that every regular tournament is arc-pancyclic (which was first proved by Alspach), and that every 2-strong tournament contains 2 distinct vertices, such that all arcs out of them are arc-pancyclic. We conjecture that there in fact exists 3 such vertices, which would be best possible (even if we looked at k -strong tournaments for any fixed $k > 1$).