

# Edge Colouring of Multigraphs

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There are two trivial lower bounds for the *chromatic index*  $\chi'(G)$  of a (multi)graph  $G$ , namely the *maximum degree*  $\Delta(G)$  and the *density*  $w(G)$ ; the last graph parameter is defined by

$$w(G) = \max_{H \subseteq G, |V(H)| \geq 2} \left\lceil \frac{|E(H)|}{\lfloor \frac{1}{2}|V(H)| \rfloor} \right\rceil.$$

A famous conjecture made, independently, by Goldberg, Anderson and Seymour in the 1970s says that every graph  $G$  satisfies

$$\chi'(G) \leq \max\{\Delta(G) + 1, w(G)\}. \quad (1)$$

If  $\chi'_f(G)$  denotes the *fractional chromatic index* of  $G$ , then (1) implies that every graph  $G$  satisfies

$$\chi'_f(G) \leq \chi'(G) \leq \chi'_f(G) + 1. \quad (2)$$

In 1990 Nishizeki and Kashiwagi proved that  $\chi'(G) \leq \max\{(11\Delta(G) + 8)/10, w(G)\}$  for every graph  $G$ . The proof was based on the so-called critical chain method. A shorter proof of this result was given by Tashkinov in 2000. The main tool in Tashkinov's proof are Tashkinov trees, a common generalization of both Vizing fans and Kierstaed paths. Favrholt, Stiebitz and Toft extended Tashkinov's method and proved in 2006 that  $\chi'(G) \leq \max\{(13\Delta(G) + 10)/12, w(G)\}$  for every graph  $G$ . In 2007 Scheide extended this result to  $\chi'(G) \leq \max\{(15\Delta(G) + 12)/14, w(G)\}$  for all graphs  $G$ . Furthermore, he proved that every graph  $G$  satisfy  $\chi'(G) \leq \max\{\Delta(G) + \sqrt{\Delta(G)/2}, w(G)\}$  and  $\chi'(G) \leq \chi'_f(G) + \sqrt{\chi'_f(G)/2}$ . The last result extends a result of Kahn from 1996 as well as a result of Sanders and Steurer from 2005. The proofs of all these results are constructive and based on an extension of Tashkinov's method. In particular, the proof of the inequality  $\chi'(G) \leq \max\{\Delta(G) + \sqrt{\Delta(G)/2}, w(G)\} =: \tau(G)$  yields an algorithm that computes, for every graph  $G = (V, E)$ , an edge colouring of  $G$  using at most  $\tau(G)$  colours, where the algorithm has time complexity bounded from above by a polynomial in  $|V|$  and  $|E|$  (and also in  $\Delta$ ).