

Assignment No. 2

Parallel Computing, DM8XX (Fall 2008)

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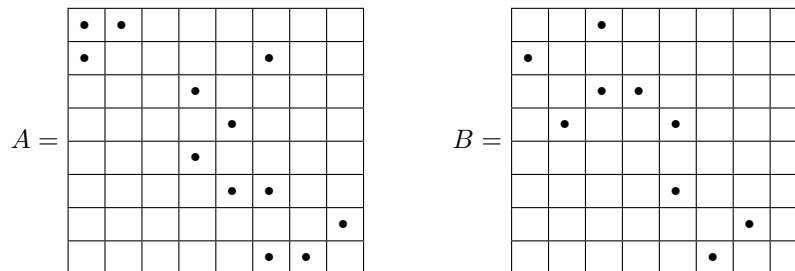
Due on: Wednesday 1. October, 12:00 p.m. (Department secretaries office (Lone Seidler Petterson) or my office).

Exercises or parts of exercises marked with * are voluntary exercises.

For preparing solutions for the following exercises you need to read the complete Chapter 3 and Sections 4.1 and 4.2 from the course book.

Exercise 1

Task Dependency Graph (5+5=10 points)



Given are the two sparse matrices A and B . Consider the problem of sparse matrix-matrix multiplication. A dot corresponds to a non-zero entry. The computation is decomposed into 8 tasks. Let task i the owner of row $A[i, *]$ and of row $B[i, *]$. Task i has to compute row i of the result $C = A \cdot B$.

- Draw the task interaction graph using directed edges. Draw an edge from task T_i to task T_j , if T_i requires data from T_j .
- Suppose that task i owns column i of matrix B instead of row i for the computation. Draw the task-interaction graph for this case.
- * Which decomposition should (usually) to be preferred? Explain why.

Exercise 2

LU factorization (5+3+3+3+3+2+2+4=25 points)

$$\begin{pmatrix} A_{1,1} & A_{1,2} & A_{1,3} \\ A_{2,1} & A_{2,2} & A_{2,3} \\ A_{3,1} & A_{3,2} & A_{3,3} \end{pmatrix} \rightarrow \begin{pmatrix} L_{1,1} & 0 & 0 \\ L_{2,1} & L_{2,2} & 0 \\ L_{3,1} & L_{3,2} & L_{3,3} \end{pmatrix} \cdot \begin{pmatrix} U_{1,1} & U_{1,2} & U_{1,3} \\ 0 & U_{2,2} & U_{2,3} \\ 0 & 0 & U_{3,3} \end{pmatrix}$$

1: $A_{1,1} \rightarrow L_{1,1}U_{1,1}$	6: $A_{2,2} = A_{2,2} - L_{2,1}U_{1,2}$	11: $L_{3,2} = A_{3,2}U_{2,2}^{-1}$
2: $L_{2,1} = A_{2,1}U_{1,1}^{-1}$	7: $A_{3,2} = A_{3,2} - L_{3,1}U_{1,2}$	12: $U_{2,3} = L_{2,2}^{-1}A_{2,3}$
3: $L_{3,1} = A_{3,1}U_{1,1}^{-1}$	8: $A_{2,3} = A_{2,3} - L_{2,1}U_{1,3}$	13: $A_{3,3} = A_{3,3} - L_{3,2}U_{2,3}$
4: $U_{1,2} = L_{1,1}^{-1}A_{1,2}$	9: $A_{3,3} = A_{3,3} - L_{3,1}U_{1,3}$	14: $A_{3,3} \rightarrow L_{3,3}U_{3,3}$
5: $U_{1,3} = L_{1,1}^{-1}A_{1,3}$	10: $A_{2,2} \rightarrow L_{2,2}U_{2,2}$	

Given is the decomposition of the LU factorization into 14 tasks. (We assume that each of the 14 tasks requires the same unit amount of work).

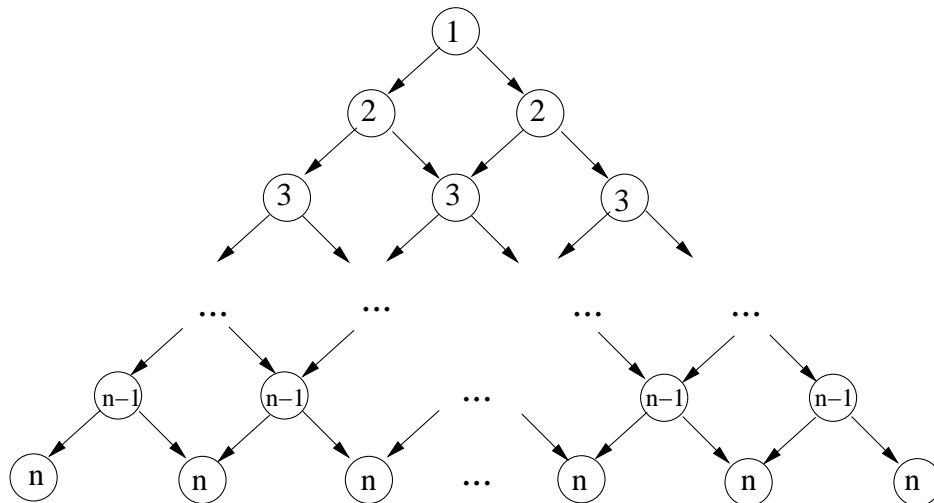
- Draw the task dependency graph.
- Determine all critical paths.
- Determine the average and the maximal degree of concurrency.
- Describe/draw an efficient mapping of the task-dependency graph of the decomposition onto three processes.
- Describe/draw an efficient mapping of the task-dependency graph of the decomposition onto four processes.
- Which of the both mappings solves the problem faster?

- g) What is the maximal speedup that can be achieved and how many processes are necessary for that speedup?
- h) What is the maximal efficiency, that can be achieved, if $p > 1$ processes are used? Describe/draw the mapping that you used.

Exercise 3

Task Dependency Graph (2+2+3+3=10 points)

Given is the following task dependency graph:



- a) Determine the maximal degree of concurrency.
- b) What is the length of the critical path?
- c) Determine the average degree of concurrency.
- d) What is the maximal speedup that can be achieved, and what is the corresponding efficiency when this speedup is realized?

Exercise 4

All-to-All Broadcast (5+5+5=15 points)

On a ring, all-to-all broadcast can be implemented in two different ways: (i) the standard ring algorithm as shown in Figure 4.9 in the course book, and (ii) the hypercube algorithm as shown in Figure 4.11. in the course book.

- a) What is the run time for case (i)?
- b) What is the run time for case (ii)?

If k messages have to traverse the same link at the same time, then assume that the effective per-word-transfer time for these messages is $k \cdot t_w$. Also assume that $t_s = 100 \cdot t_w$.

- c) Which of the two methods, (i) or (ii), is better if the message size m is very large? Which method is better if m is very small (may be one word)? Explain.

Exercise 5*

All-to-All Broadcast on a Tree (voluntary, but not very complicated)

Given a balanced binary tree as shown in Figure 4.7 from the course book, describe a procedure to perform all-to-all broadcast that takes time $(t_s + t_w \cdot m \cdot p/2) \log p$ for m -word messages on p nodes. Assume that only the leaves of the tree contain nodes, and that an exchange of two m -word messages between any two nodes connected by bidirectional channels takes time $t_s + t_w \cdot m \cdot k$ if the communication channel (or a part of it) is shared by k simultaneous messages.