

Problem 1 Given a FA  $M$   
and  $w \in \Sigma^*$  is  $w \in L(M)$ ?

Problem 2 Given FA  $M$   
is  $L(M) = \emptyset$ ?

Problem 3 Given FA's  $M_1, M_2$ , is  $L(M_1) \subseteq L(M_2)$ ?

Problem 4 Given FA's  $M_1, M_2$ , is  $L(M_1) = L(M_2)$ ?

Problem 5 Given FA  $M$ , is  $L(M) = \Sigma^*$ ?

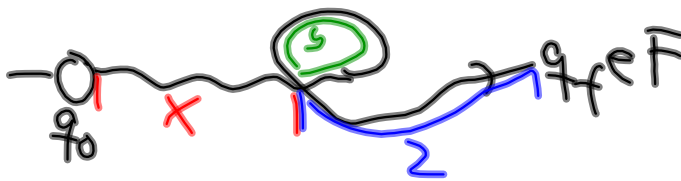
Problem 6 Given regular expr.  $\alpha$ , construct  $M$  s.t.  $L(\alpha) = L(M)$

Problem 7 Given NFA  $M$ , Find reg. exp.  $\alpha$ , s.t.  $L(M) = L(\alpha)$ .

$L$  regulært  $\Leftrightarrow \exists$  DFA  $M$   
 så  $L = L(M)$

Lad  $q = \#$  tilst. i  $M$ .

se på  $w$  så  $|w| \geq p$



$w = xyz$   
 $|xy| \leq p$   
 $|y| > 0$

alternativ form:

1.  $w = \bar{x}y\bar{z}$
2.  $|\bar{y}z| \leq p$
3.  $|\bar{y}| > 0$

Ex 1.77  $E = \{0^i 1^j \mid i > j\}$

anvend a lternative version (baglæns) af PL

$$S = 0^{p+r} 1^p \text{ hvor } p \text{ er pumpelængd.}$$

$$r \geq 1$$

Modstander returnerer  $x, y, z$  så

$$0^{p+r} 1^p = xyz, \quad |yz| \leq p, \quad |y| > 0$$

se på  $xy^{|y|}z$ :  $0^{p+r} 1^{p+r \cdot |y|} \notin E \rightarrow \times$

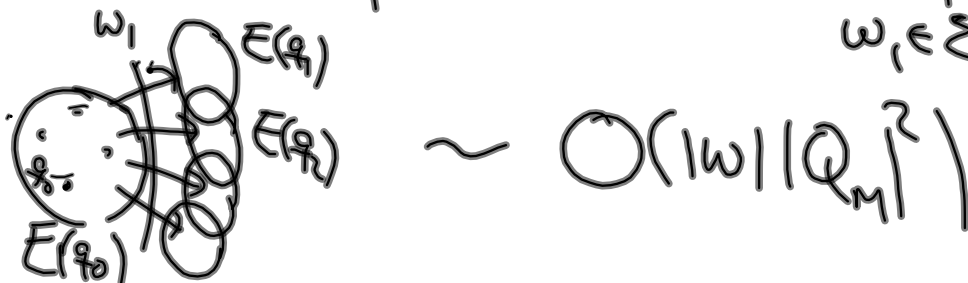
Problem 1: • Hvis  $M$  er en DFA

Så kør  $M$  på  $w$ :  $w \in L \Leftrightarrow M$  er i en accept state.

• Hvis  $M$  er en NFA

– Lav ækv. DFA i exp tid  $2^{|Q|}$ . pol  
og check som ovenfor.

– Simuler  $M$  på  $w$ :  $w = w_1 w_2 \dots w_n$   
 $w_i \in \Sigma$



Problem 2 Uanset om  $M$  er DFA  
eller NFA:

$$L(M) \neq \emptyset$$



$$\exists (q_0, F)\text{-vej i } M.$$



Problem 3

$$L(M_1) \subseteq L(M_2)$$
$$\iff L(M_1) \cap \overline{L(M_2)} = \emptyset$$

$= L(M)$   
for an FA  $M$ .  
by 2.

$$S \rightarrow aSa | bSb | A, \quad A \rightarrow a | b | \epsilon$$

1.  $S_0 \rightarrow S$  new variable  $S_0$

2. order variables  $A, S$

$$S \rightarrow aSa | bSb | \epsilon, \quad A \rightarrow a | b \quad (\text{Eliminate } A \rightarrow \epsilon)$$

$$S \rightarrow aa | aSa | bb | bSb \quad (\text{Eliminate } S \rightarrow \epsilon)$$

3.  $S_0 \rightarrow aa | aSa | bb | bSb$   
 $S \rightarrow aa | aSa | bb | bSb$   
 $A \rightarrow a | b$       Eliminate  $S_0 \rightarrow S$

4.  $S_0 \rightarrow aa | aV_{Sa} | bb | bV_{Sb}$   
 $S \rightarrow aa | aU_{Sa} | bb | bU_{Sb}$   
 $V_{Sa} \rightarrow Sa, V_{Sb} \rightarrow Sb$   
 $U_{Sa} \rightarrow Sa, U_{Sb} \rightarrow Sb$        $A \rightarrow a | b$

Final (reduced)

$$S_0 \rightarrow AV_{S_a} | BV_{S_b} | AA | BB$$

$$S \rightarrow AV_{S_a} | BV_{S_b} | AA | BB$$

$$V_{S_a} \rightarrow SA$$

$$V_{S_b} \rightarrow SB \quad + \quad \underline{S \rightarrow \epsilon}$$

$$A \rightarrow a$$

$$B \rightarrow b$$

$$R: \begin{cases} A \rightarrow 0A1 & | B \\ B \rightarrow \# \end{cases} \quad G_1$$

$$A \Rightarrow 0A1 \Rightarrow 00A11 \Rightarrow 00B11 \\ \Rightarrow 00\#11$$

$$G = (V, \Sigma, R, S)$$

$$L(G) = \{ w \mid S \xRightarrow{*} w \}$$

$$L(G_1) = \{ 0^n \# 1^n \mid n \geq 0 \}$$

Hvis  $L$  er regulært så er

$L = L(G)$  for en CFG  $G$ .

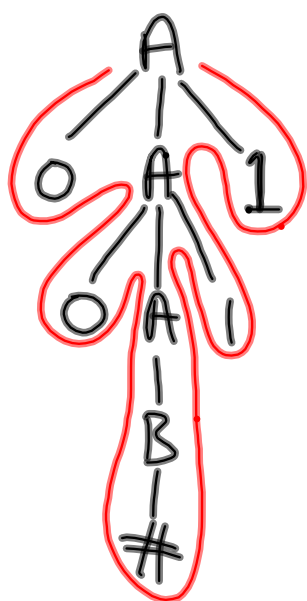
P: Lad  $M = (Q, \Sigma, \delta, q_0, F)$  DFA:  $L(M) = L$

Sæt  $V = Q$

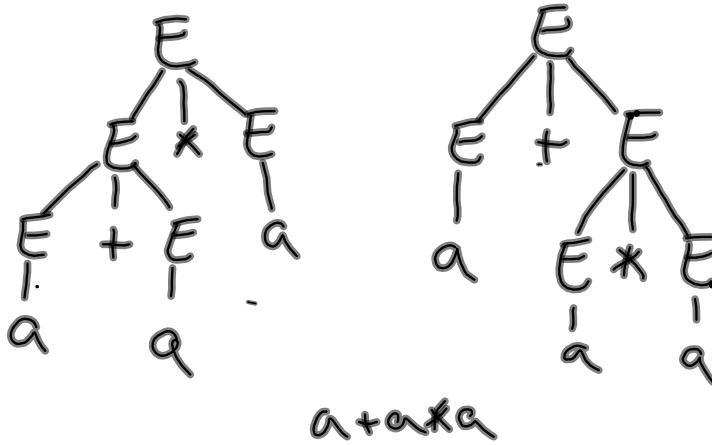
$S = q_0$

$R: \{ q_i \rightarrow a q_j \mid \delta(q_i, a) = q_j \}$

$\cup \{ q_f \rightarrow \epsilon \mid q_f \in F \}$



$G_5 \quad E \rightarrow E + E \mid E * E \mid (E) \mid a$



Chomsky normal form.

$$R: A \rightarrow BC \quad B, C \in V - S$$

$$A \rightarrow a \quad a \in \Sigma$$

$$+ S \rightarrow \epsilon$$

Thm

$$\forall \text{ CFG } G$$

$$\exists \text{ CNF CFG } G' : L(G) = L(G')$$

1. Tilføj  $S_0 \rightarrow S$   $S_0$  ny var.

2. Fjern  $\epsilon$  produktioner ( $C \rightarrow \epsilon$ )

$$R \rightarrow uAvAw$$

$$\downarrow R \rightarrow uAvAw \mid u\epsilon Aw \mid uA\epsilon w$$

?  $B \rightarrow A$  ( $A \rightarrow \epsilon$ )

$$B \rightarrow \epsilon$$

3.  $R \rightarrow A$  ( $A \rightarrow u$ )

$$\downarrow R \rightarrow u$$

$$A \rightarrow u_1 u_2 \dots u_k \quad k \geq 3 \quad u_i \in V \cup \Sigma$$

$$\downarrow$$

$$\begin{array}{l} A \rightarrow u_1 u_2 \\ u_2 \rightarrow u_2 u_3 \\ \vdots \\ u_{k-2} \rightarrow u_{k-1} u_k \end{array} \quad \left| \quad \begin{array}{l} ? \quad u_3 \rightarrow a u_4 \\ \downarrow \\ u_3 \rightarrow A u_4 \\ A \rightarrow a \end{array}$$