

---

---

---

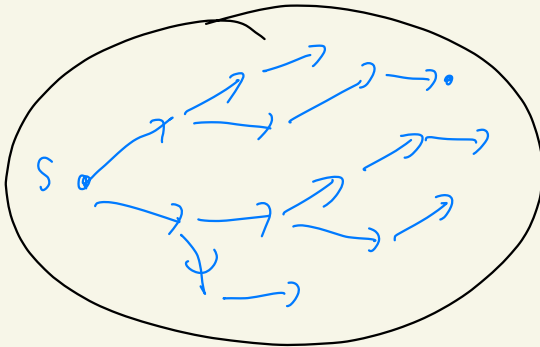
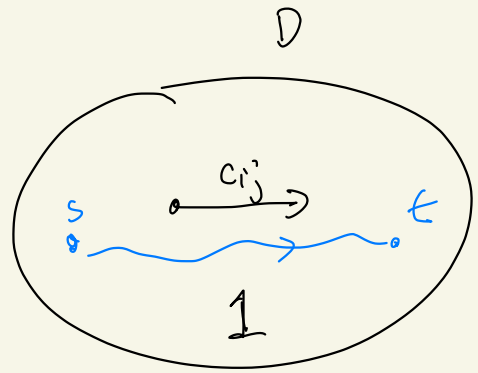
---

---



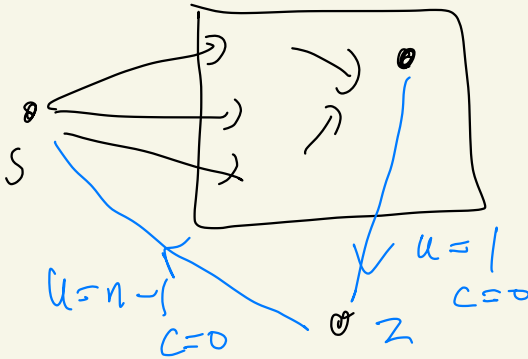
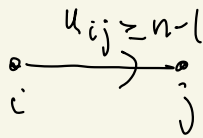
# Ahuja 1.1

Shortest path problem

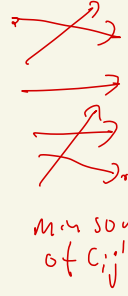
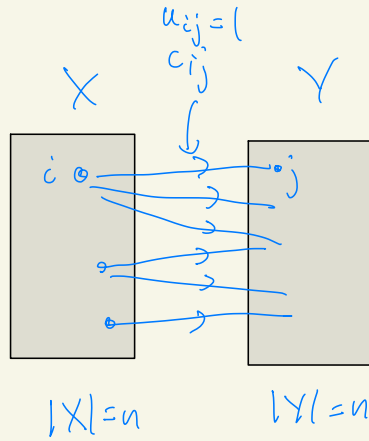


$$b_x(s) = n - 1$$

$$b_x(v) = -1 \quad \forall v \neq s$$

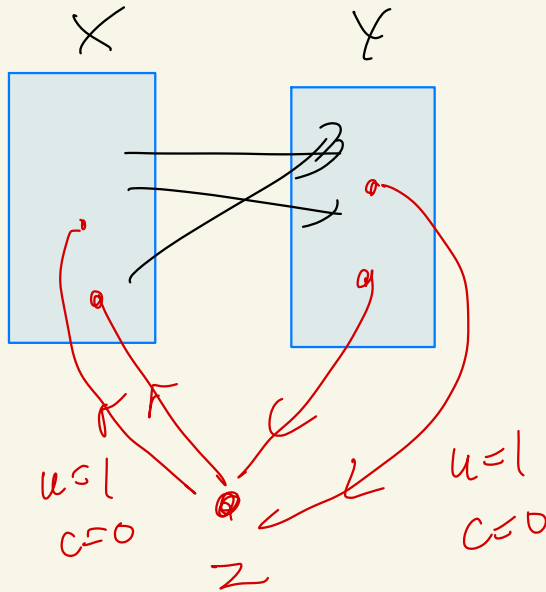


# assignment

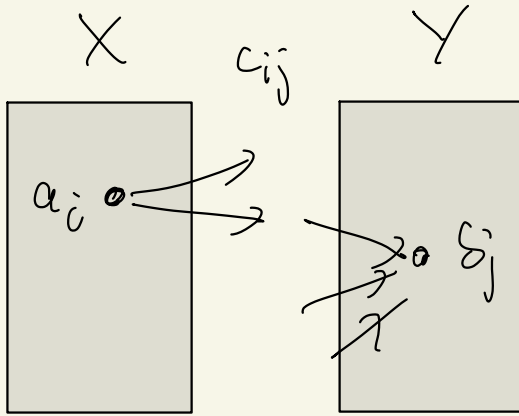


$$b_X(\sigma) = 1 \text{ if } \sigma \in X$$

$$b_X(\sigma) = -1 \text{ if } \sigma \in Y$$

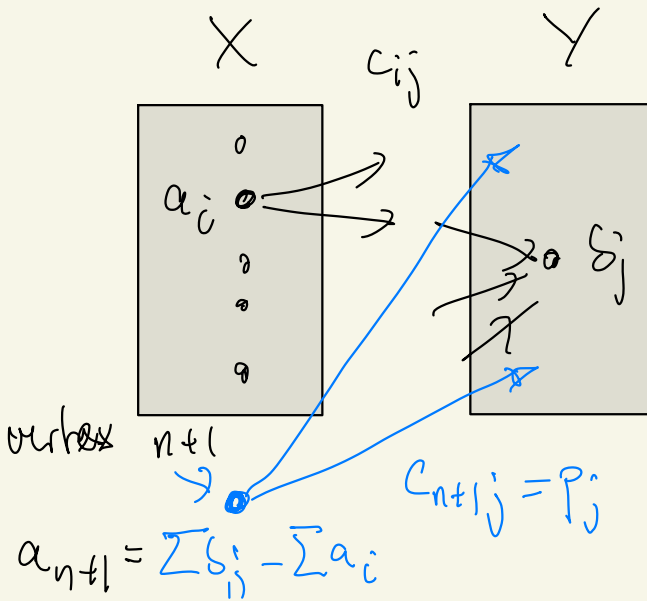


# Ahuja 1.2

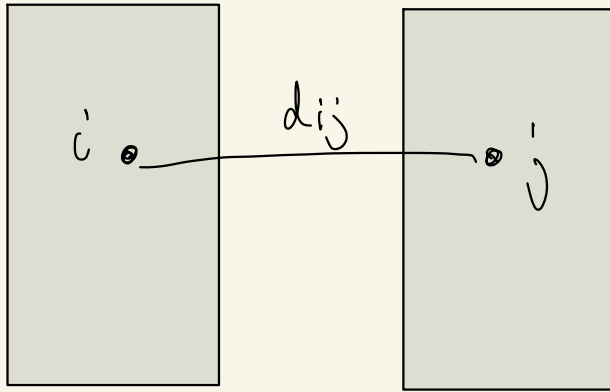


transportation problem  $\sum_i a_i = \sum_j b_j$

suppon  $\sum b_j > \sum a_i$



# Ahuja 1.4



Persons

Jobs

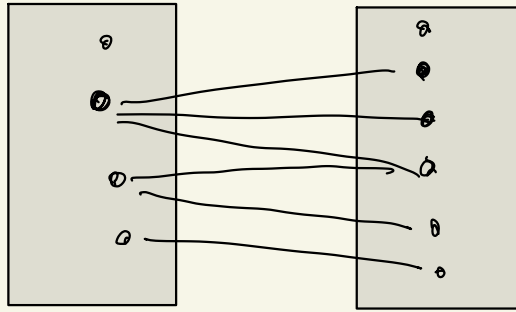
$d_{ij}$  = utility of assigning person  
i to job j

want to maximize total utility

$$\text{cost } c_{ij} = -d_{ij}$$

Ahuja 1.5

# Dating problem



$X$

$Y$

$G$

find max no of compatible  $x-y$

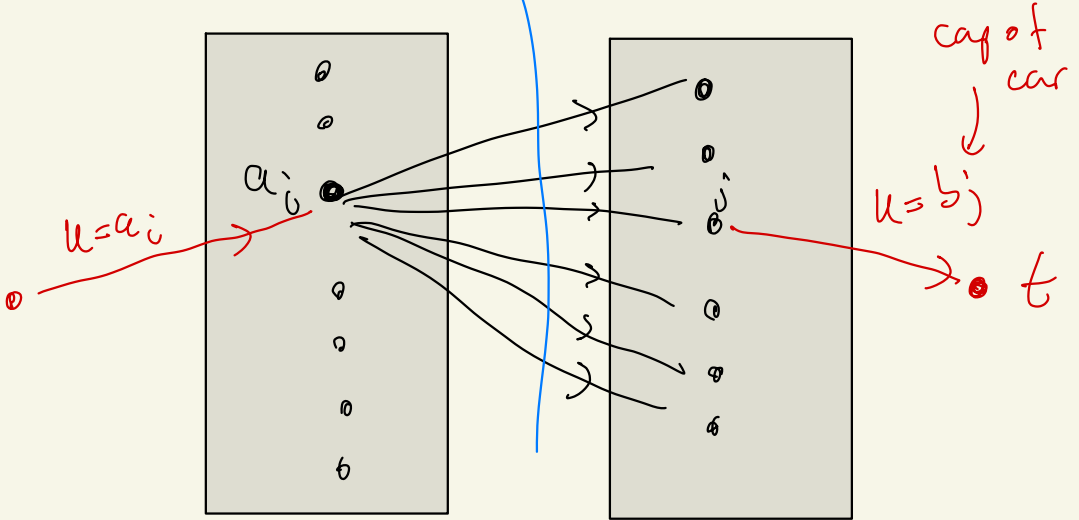
= max matchings in  $G$

# Almuj 9 1.8

F-Families

$u=1$

Cars



$$\sum_{i \in F} a_i$$

integer

good distribution  $\longleftrightarrow$  flow of value

$$\sum_{i \in F} a_i$$



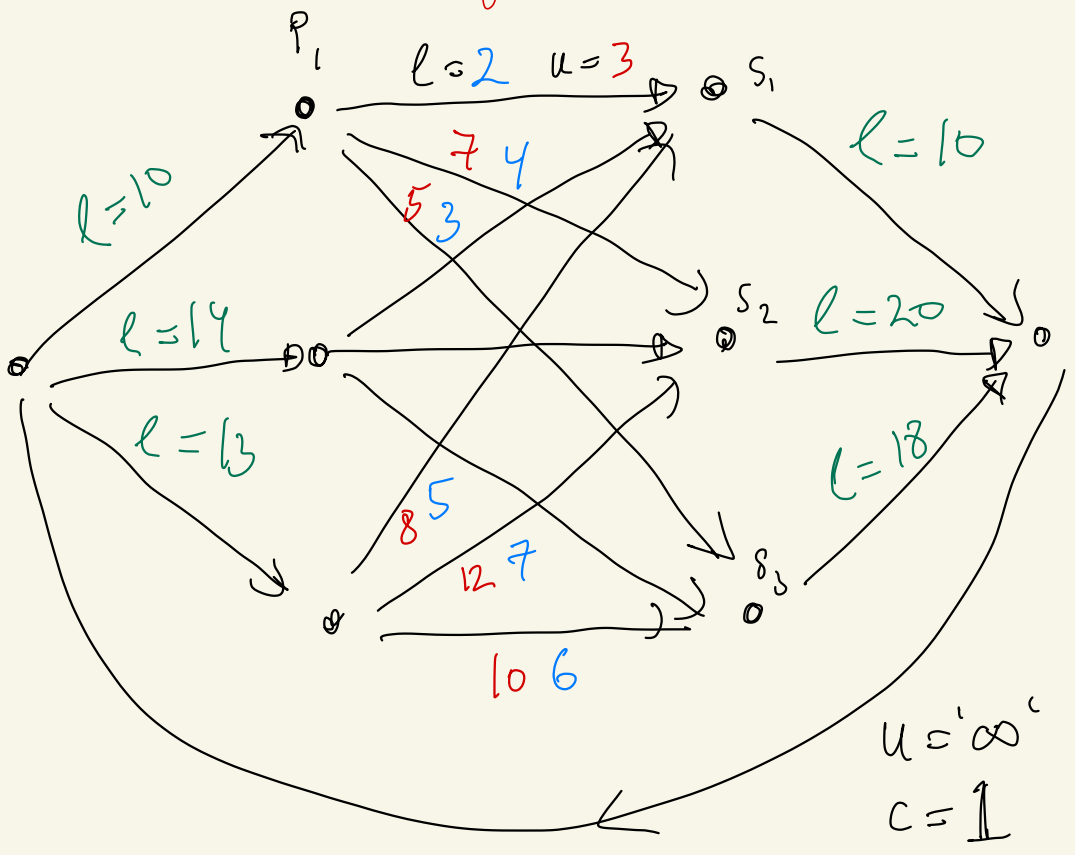


Ahuja 1.9

$\geq 10$   
 $\geq 14$   
 $\geq 13$

	$\geq 10$	$\geq 20$	$\geq 18$
	$s_1$	$s_2$	$s_3$
$p_1$	2 3	4 7	3 5
$p_2$	3 5	6 7	5 10
$p_3$	5 8	7 12	6 10

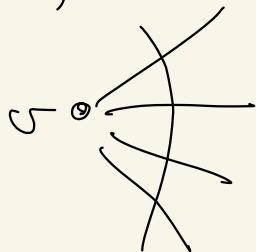
$i$  = minimum # of cars  
 $j$  = maximum # of cars



Ahuja 2.12

$$G = (V, E)$$

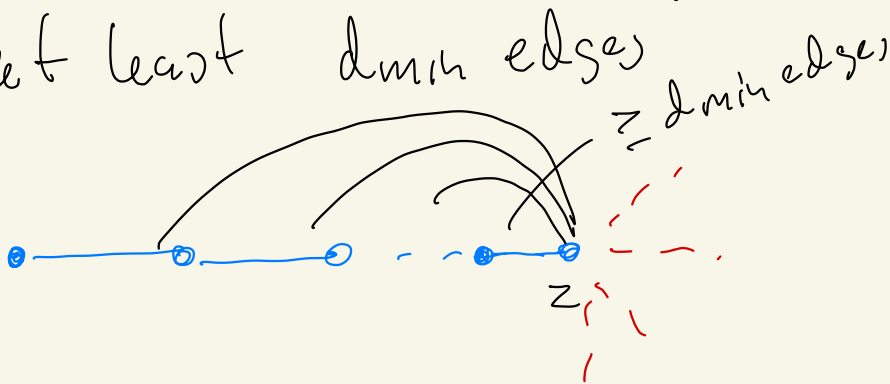
$d(v)$



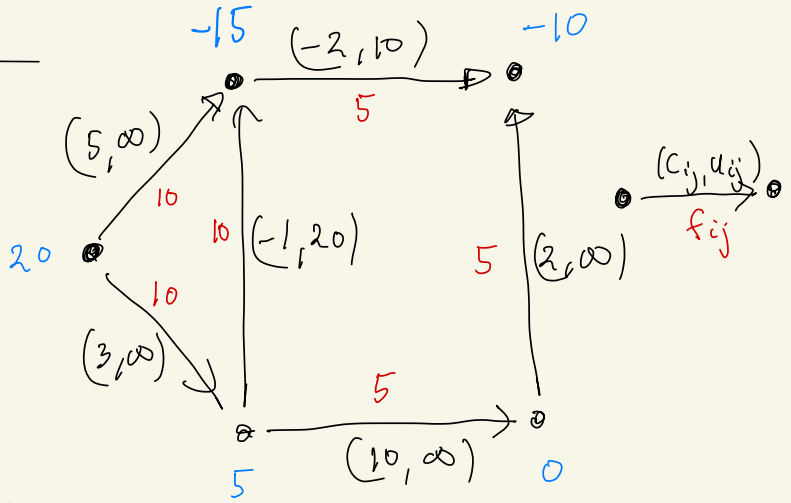
# of edges  
incident to v

$$d(v) \geq d_{\min} \quad \forall v$$

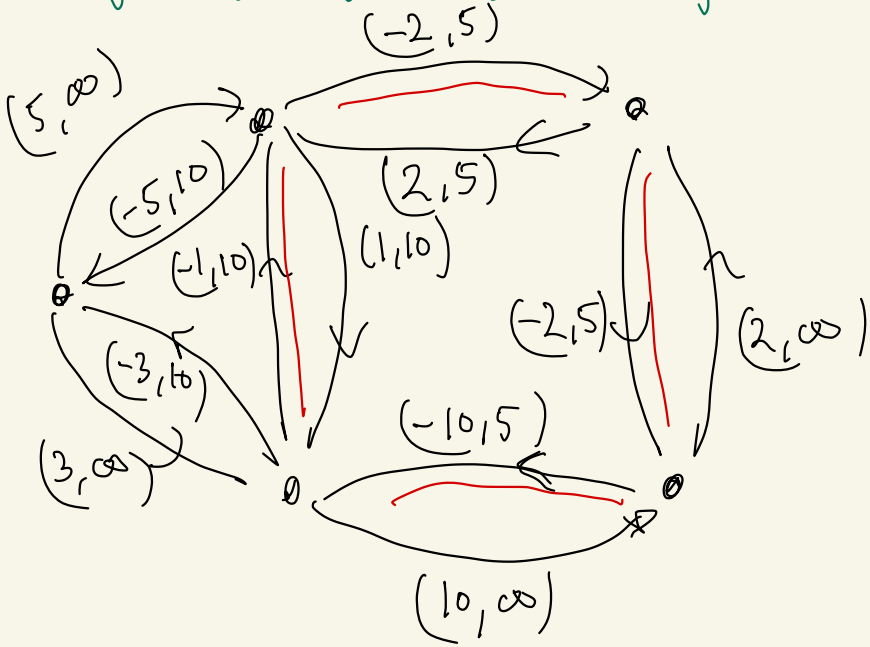
Show that  $G$  has a path with  
at least  $d_{\min}$  edges



# Alwja 2.45



$$r_{ij} = (u_{ij} - x_{ij}) + x_{ij} \quad l_{ij} = 0 \quad \forall ij$$

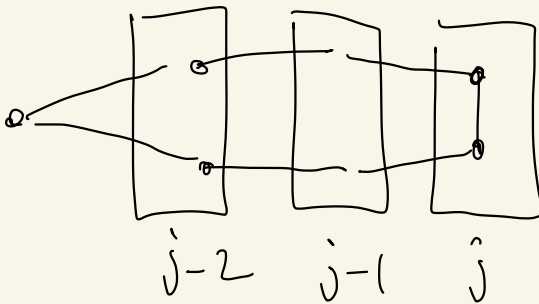
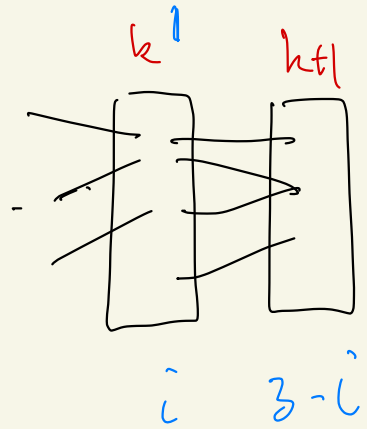
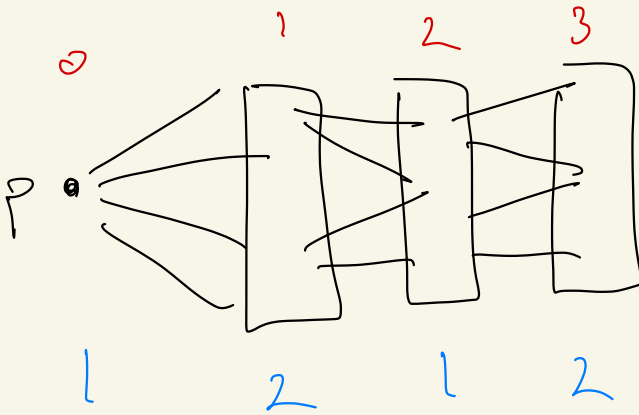
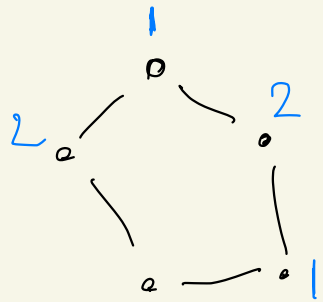
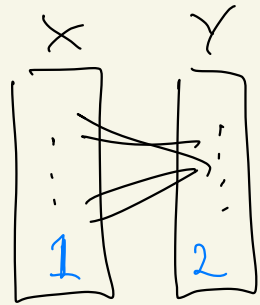


red cycle has cost  $-15$   
 so flow is not min cost

$G$  is bipartite



$G$  has no cycle of odd length



edge inside distance class  $\Leftrightarrow$  odd cycle

prove that a strongly connected digraph is bipartite if and only if it has no directed cycle of odd length

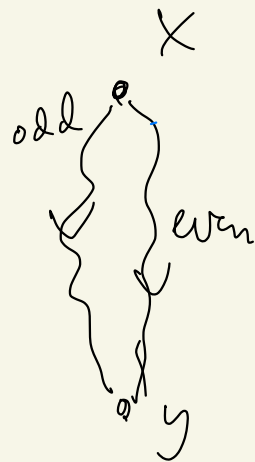
$\Rightarrow$  ✓

$\Leftarrow$  fix vertex  $x$

claim:  $\forall y \in V - x$

$$|A(P_1)| \equiv |A(P_2)| \pmod{2}$$

$\forall (x, y)$ -paths



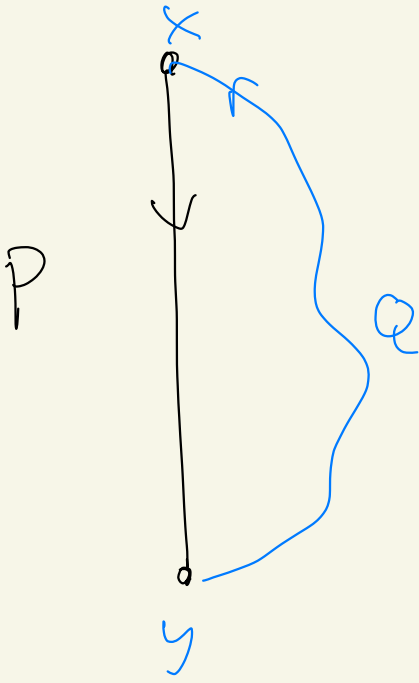
proof suppose not. Choose  $P$  and  $Q$



s.t.  $|A(P)| \not\equiv |A(Q)| \pmod{2}$

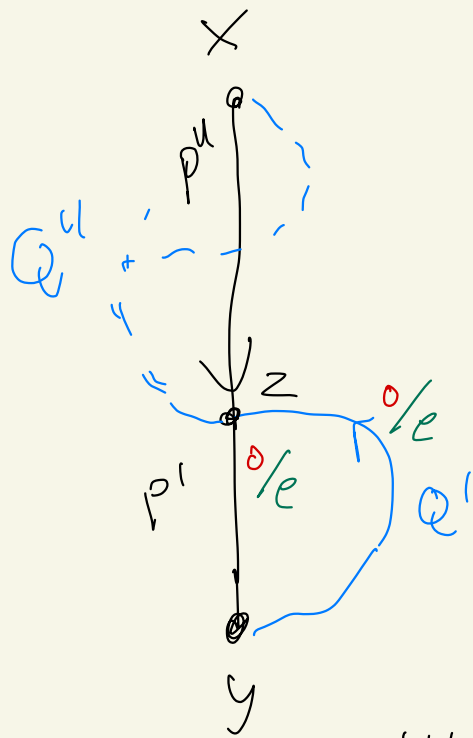
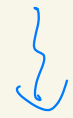
and  $|A(P)| + |A(Q)|$  minimum

(\*)



if  $V(P) \cap V(Q) = \{x, y\}$

odd cycle



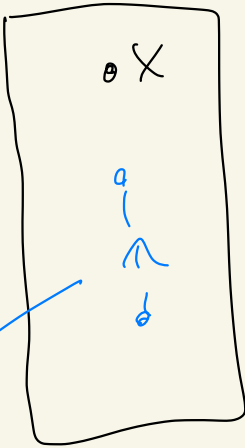
by  $\otimes$   $|A(P')| \equiv |A(Q')| \pmod{2}$

$\Rightarrow |A(P'')| \not\equiv |A(Q'')| \pmod{2} \downarrow$

define

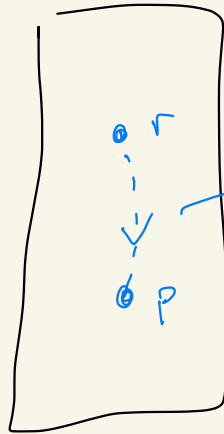
$U_1 = \{y \mid \text{Every } (x,y)\text{-path is even}\}$

$U_2 = \{y \mid \text{Every } (x,y)\text{-path is odd}\}$



no arc

$U_1$

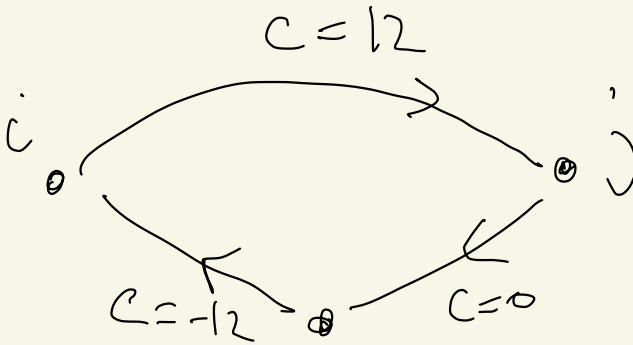
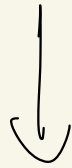
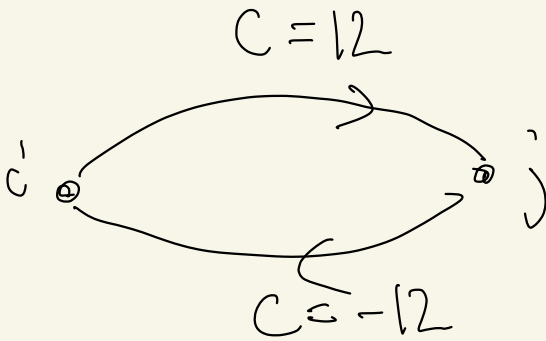


no arc

$U_2$

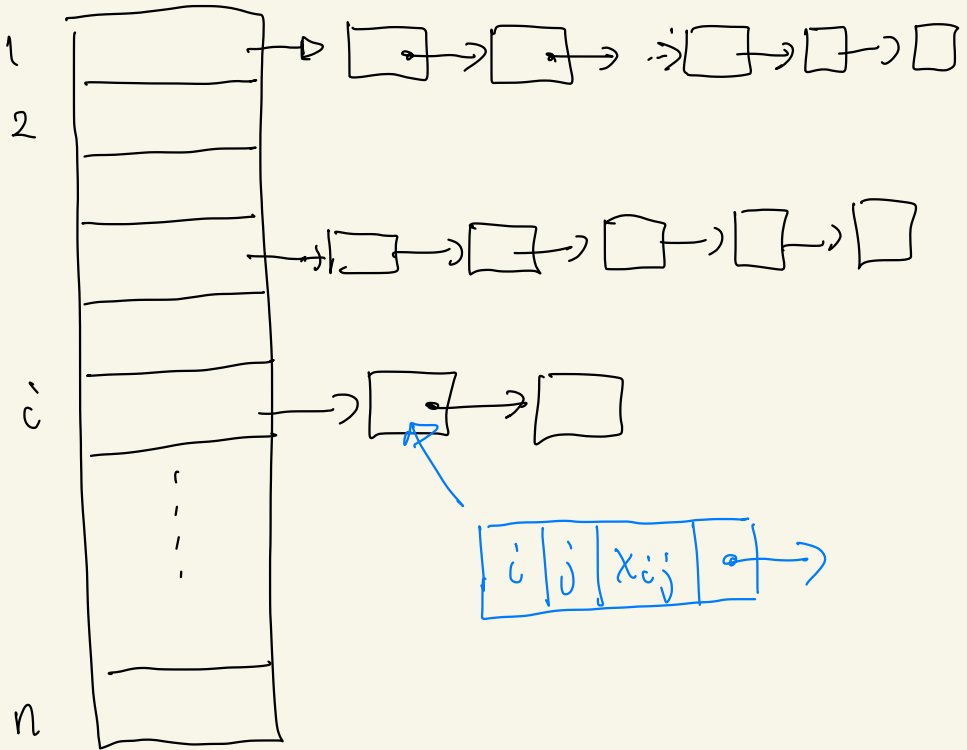
$\Rightarrow D$  is bipartite.

# BJG 3.2

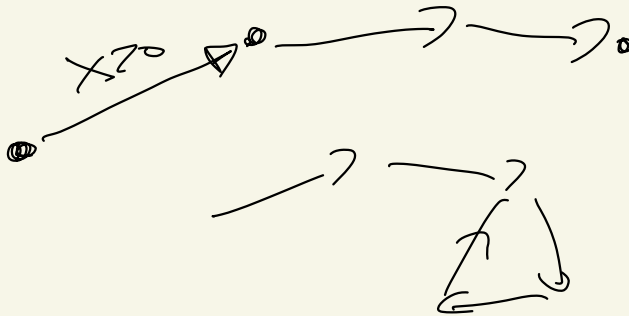




# BJG 3.7



adjacency lists



# B)G 3.8

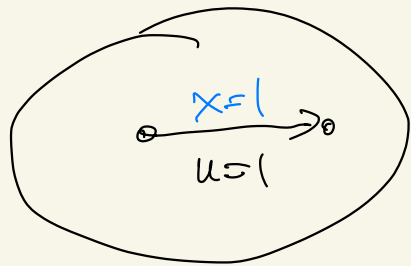
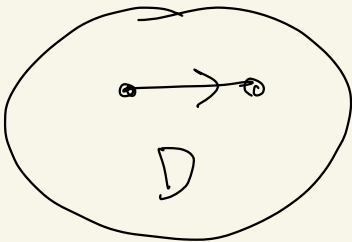
---

$D$  is eulerian

$\Downarrow$  def

$$d^+(v) = d^-(v) \quad \forall v$$

Show  $A(D)$  (arc set of  $D$ ) can be decomposed into arc-disjoint cycles



$$N_D = (V, A, \ell=0, \varphi=1)$$

$x$  is a circulation

$\Rightarrow x$  decomposes in cycleflows

$\Rightarrow$  arc-disj cycles  $w_1, w_2, \dots, w_r$