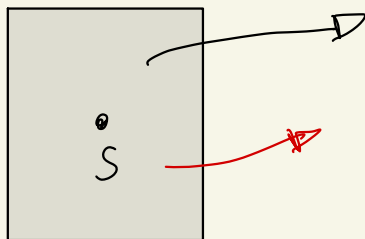
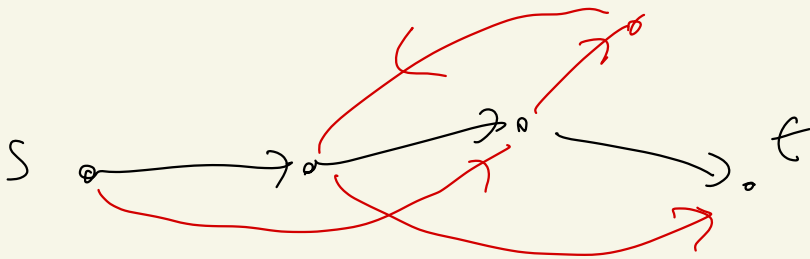



Combinatorial application
of the max flow min cut
theorem

The arc-connectivity from
 s to t in a digraph $D=(V,A)$
with $s,t \in V$, called $\lambda(s,t)$
is the maximum number of
arc-disjoint (s,t) -paths in D



S

t

$$d^+(s) \geq \lambda(s,t)$$

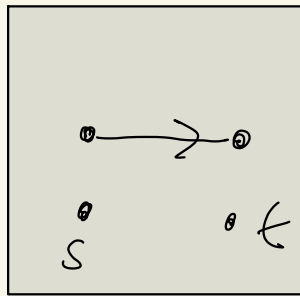
Theorem (Menger)

Let $D=(V, A)$ be a digraph and

let $s, t \in V$. Then

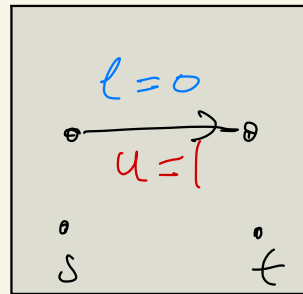
$$\lambda(s, t) = \min \{ d^+(s) \mid s \in S, t \in V-S \}$$

P:



$$V = V' \cup \{s, t\}$$

$$D = (V, A)$$



$$N(D)$$

$$N(D) = \{ V' \cup \{s, t\}, A, l, u \}$$

Claim: The maximum value of an (s, t) -flow in $N(D)$ is $\lambda(s, t)$

$$\max |x| \geq \lambda(s, t):$$

Given P_1, P_2, \dots, P_k arc-disj
 (s, t) -paths in D

send 1 unit of flow along

each P_i in $N(D) \rightarrow$ gives a

feasible (s, t) -flow x with $|x| = k$

$$\lambda(s, t) \geq |x^*|: \quad x^* \text{ max flow}$$

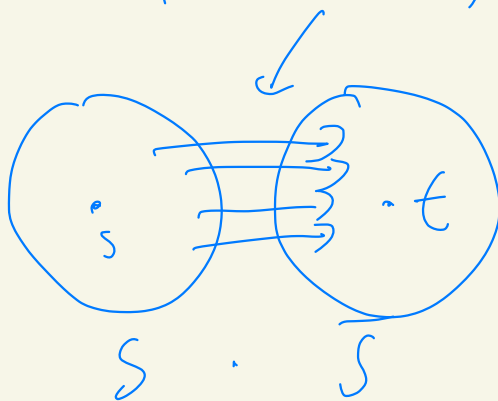
Given an integer valued
maximum flow x^*

Decompose x^* into path flows and
cycle flows $f(P_1^i) \dots f(P_{|x^*|}^i), f(C_1) \dots f(C_r)$

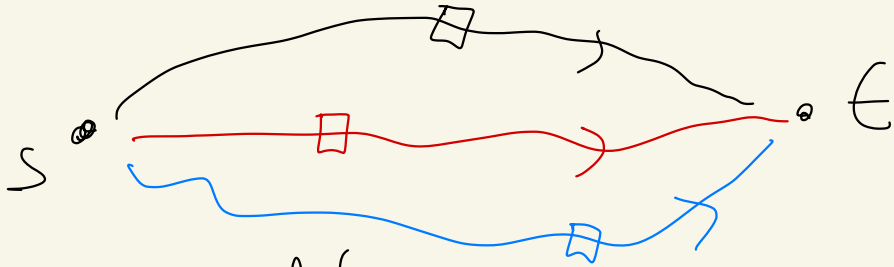
each P_i^i is an (s, t) -flow

Each P_i^1 is a path
(the same in D) and
they are arc-disjoint
(since $u_{ij} = 1 \quad \forall i, j \in A$)

$$\lambda(s, t) = |X^*| = u(s, \bar{s}) = \underline{d^t(s)}$$



Internally disjoint (s,t) -paths
in digraphs



def
 $k(s,t) = \max \#$ of internally
disjoint (s,t) -paths

assume $s \neq t$

(s,t) -separator is a set $X \subseteq V - \{s,t\}$

s.t. $D - X$ has no (s,t) -path

(that is, every (s,t) -path uses at
least one vertex in X)

Claim: $k(s,t) \leq \min \{ |X| \mid X \text{ is an } (s,t)\text{-sep} \}$

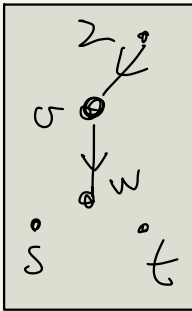
Theorem (Menger) $D = (V, A)$

$$s, t \in V$$

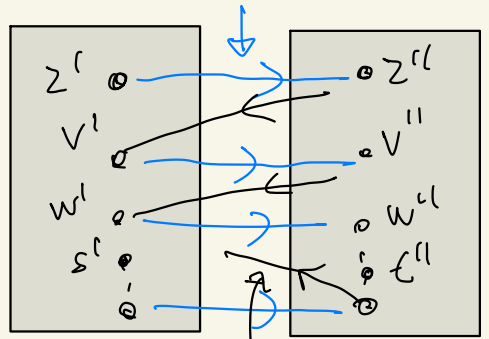
So upon $st \notin A$

Then $k(s, t) = \min \{ |X| \mid X(s, t) = \emptyset \}$

$$V^r \quad u=1 \quad V^l$$



D



$N_D \quad u=\infty$

claim

$k(s, t) = \max$ value of an (s'', t') -flow in N_D

claim

$K(s,t) = \max$ value of an
 (s'',t') -flow in N_D

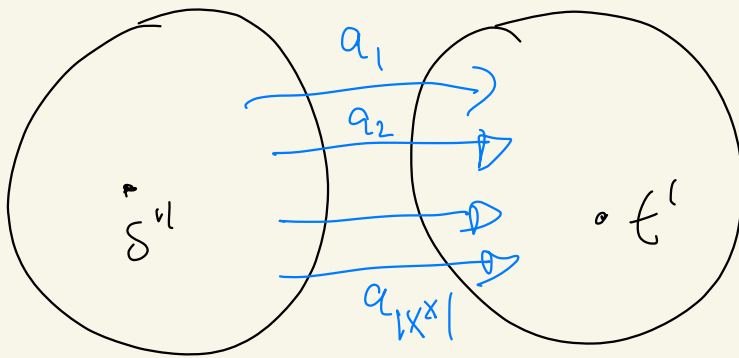
$P:$ \leq let P_1, \dots, P_g be internally
disjoint (s,t) -paths in D
and let P'_1, \dots, P'_g be the
corresponding paths in N_D
send 1 unit along each P'_i .

\geq let x^* be a max (s'',t') -flow
(integer valued)
 x^* can be decomposed into $|x^*|$
path flows $f(Q_1), \dots, f(Q_{|x^*|})$
and some cycle flows
each Q_i, Q_j are internally
disj.

So $K(s,t) \geq |x^*|$

$$k(s, t) = |X^*| = u(s, \bar{s})$$

$$= \# \text{ shortest paths from } s \text{ to } \bar{s}$$



each $a_i \leftrightarrow$ a vertex v_i in D
 and removing $v_1, v_2, \dots, v_{|X^*|}$
 kills all (s, t) -paths in D
 so $X = \{v_1, v_2, \dots, v_{|X^*|}\}$ is an
 (s, t) -separator with $k(s, t) = |X| \quad \square$