

## Exercise 8.5

We will use Gomory's cutting plane algorithm. We add a surplus variable  $s_4$  to the constraint and multiply the objective by  $-1$ , to obtain a maximization problem. The simplex tableau for the problem is then:

max	$x_1$	$x_2$	$x_3$	$s_4$	
-1	5	9	23		0
	-20	-35	-95	1	-319

All coefficients of the objective are positive and the right hand side of the constraint is negative, hence the solution is dual feasible but not optimal. Thus we can use dual-simplex to find the optimal solution. We do the coefficient test (row marked "Test") and find the least negative number ( $\frac{23}{-95}$ ).

max	$x_1$	$x_2$	$x_3$	$s_4$	
-1	5	9	23		0
	-20	-35	-95	1	-319
Test	$\frac{5}{-20}$	$\frac{9}{-35}$	$\frac{23}{-95}$		

Pivoting the circled element gives the optimal simplex tableau:

max	$x_1$	$x_2$	$x_3$	$s_4$	
-1	15/95	50/95	23/95		$-7337/95 \approx -77, 23$
	20/95	35/95	1	-1/95	319/95

The right hand side of the row is not integer (and hence neither is  $x_3$ ), so we use this row to get a Gomory cut. The Gomory cut obtained is:

$$\begin{aligned} & \frac{20}{95}x_1 + \frac{35}{95}x_2 + \frac{94}{95}s_4 \geq \frac{34}{95} \\ \Rightarrow & -\frac{20}{95}x_1 - \frac{35}{95}x_2 - \frac{94}{95}s_4 + s_5 = -\frac{34}{95} \end{aligned}$$

Adding this row give us the following tableau. Again we need dual-simplex to reoptimize so the coefficient test is shown including the selected pivot element.

max	$x_1$	$x_2$	$x_3$	$s_4$	$s_5$	
-1	15/95	50/95	23/95			$-7337/95 \approx -77, 23$
	20/95	35/95	1	-1/95		319/95
	-20/95	-35/95		-94/95	1	-34/95
Test	$-\frac{15}{-20}$	$-\frac{50}{-35}$		$-\frac{23}{-94}$		

Carrying out the pivoting gives the following optimal tableau.

max	$x_1$	$x_2$	$x_3$	$s_4$	$s_5$	
-1	10/94	41/94			23/94	$-7268/94 \approx -77, 32$
	20/94	35/94	1		-1/94	316/94
	20/94	35/94		1	-95/94	34/94

We note that the original objective  $z$  is increasing ( $77, 23 \rightarrow 77, 32$ ) thus the lower bound on the solution value gets better - as expected. This process can be continued while improving the lower bound on the solution. At this particular instance any of the two rows can be selected to generate a Gomory cut, since both have right hand sides. The rate of improvement is in general, however, very slow.

## Exercise 8.8

Denote the nodes in the odd hole with  $H = \{1, 2, \dots, k\}$  ( $k$  odd) such that the edges of the odd hole are  $12, 23$  and  $k1$ . For all edges  $ij$  it holds that  $x_i + x_j \leq 1$ . Add all inequalities corresponding to edges in the odd hole and obtain:

$$\begin{aligned}
2x_1 + 2x_2 + \dots + 2x_k &\leq |H| \\
\Rightarrow x_1 + x_2 + \dots + x_k &\leq \frac{|H|}{2} = \frac{|H| - 1}{2} + \frac{1}{2} \\
\Rightarrow x_1 + x_2 + \dots + x_k &\leq \frac{|H| - 1}{2}
\end{aligned}$$

The last  $\Rightarrow$  follows from the fact that since  $|H|$  is odd  $\frac{|H|-1}{2}$  is integer and accordingly all variables on the left hand side are integers, hence  $\frac{1}{2}$  can be left out.

## Exercise 8.15

(i)

For  $u = 1/4$ , the following inequality is obtained.

$$\begin{aligned}
\frac{3}{4}x_1 - x_2 + \frac{1}{2}x_3 - \frac{3}{4}x_4 + \frac{1}{4}x_5 &\leq -\frac{1}{2} \\
\Rightarrow -x_2 - x_4 &\leq -1 \Rightarrow x_2 + x_4 \geq 1
\end{aligned}$$

Remember that implicitly we have an upper bound on the variables (binary variables) thus  $x_i \leq 1$  are valid inequalities.

Taking 1 time the original constraint and 1 time  $x_2 \leq 1$ , the following constraint is obtained:

$$3x_1 + (1 - 4)x_2 + 2x_3 - 3x_4 + x_5 \leq 1 - 2$$

Multiplying through by  $1/3$  for this constraint we get:

$$\begin{aligned}x_1 - x_2 + \frac{2}{3}x_3 - x_4 + \frac{1}{3}x_5 &\leq -\frac{1}{3} \\ \Rightarrow x_1 - x_2 - x_4 &\leq -1\end{aligned}$$

Adding  $x_4 \leq 1$  the desired inequality is obtained:

$$x_1 - x_2 + (1 - 1)x_4 \leq 1 - 1 \Rightarrow x_1 - x_2 \leq 0 \Rightarrow x_1 \leq x_2$$

(ii)

$u_{ij}$  corresponds to the constraint  $x_i + x_j \leq 1$ . Set  $u_{12} = u_{13} = u_{23} = 1/2$  and remaining  $u$ 's = 0. This gives:

$$\begin{aligned}\frac{1}{2}(x_1 + x_2) + \frac{1}{2}(x_1 + x_3) + \frac{1}{2}(x_2 + x_3) &\leq 3 \cdot \frac{1}{2} \\ \Rightarrow x_1 + x_2 + x_3 &\leq 1\end{aligned}$$

Correspondingly the following inequalities are valid  $x_1 + x_2 + x_4 \leq 1$ ,  $x_1 + x_3 + x_4 \leq 1$  and  $x_2 + x_3 + x_4 \leq 1$ . Letting  $u_{123} = u_{124} = u_{134} = u_{234} = 1/3$  and remaining  $u$ 's = 0, the desired inequality ( $x_1 + x_2 + x_3 + x_4 \leq 1$ ) is obtained.