

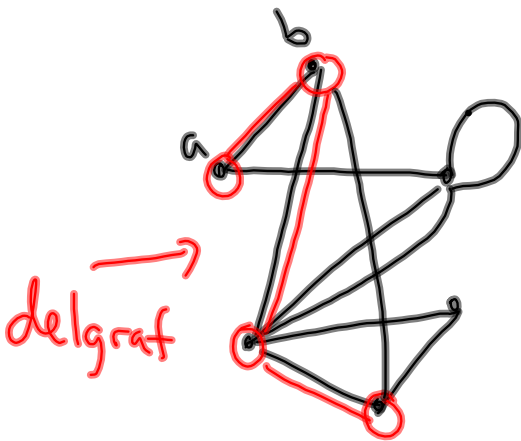
En graf  $G = (V, E)$

$V$  punkterne i  $G$   
(Knuderne)

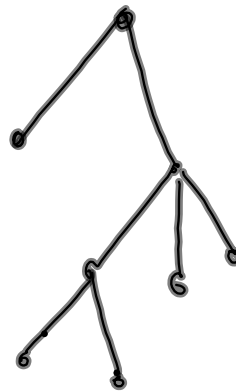
$$\frac{\text{Toft}}{P(G) = V}$$

$E$  kanten i  $G$

$$K(G) = E$$



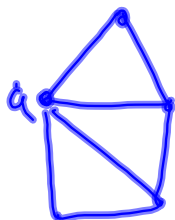
$$E \subseteq V \times V$$



tree

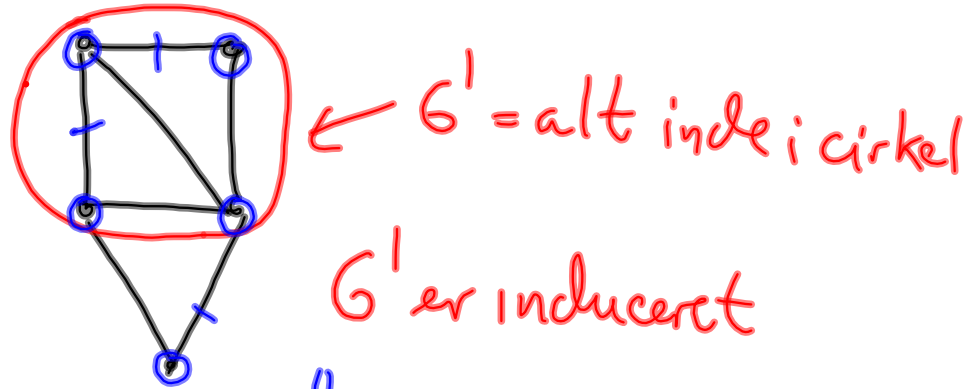
valensen (graden) af et  
punkt  $x \in V$

= # Kanter med endepunkt  
i  $x$



$$v(a, G) = 4$$

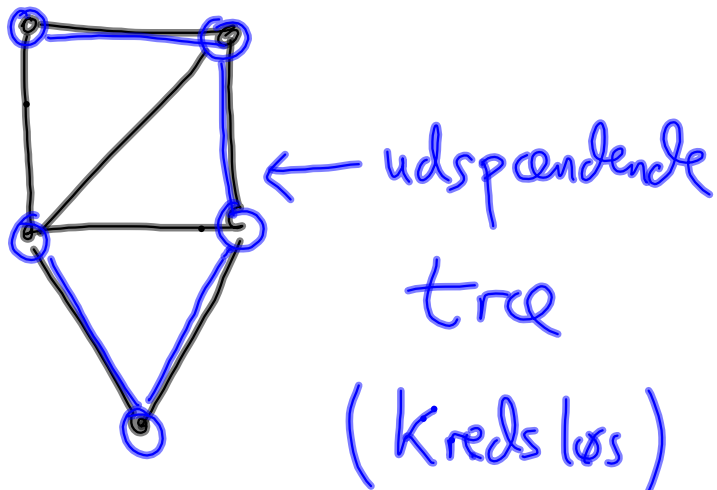
$$d(a) = 4$$

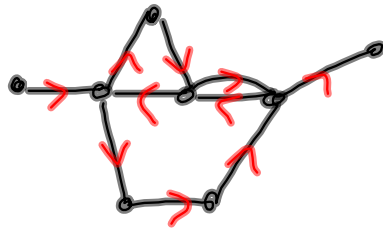


$G$   $G''$  udspondene i  $G$

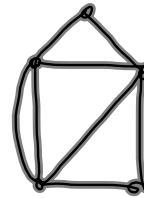
↑

$$P(G'') = P(G)$$



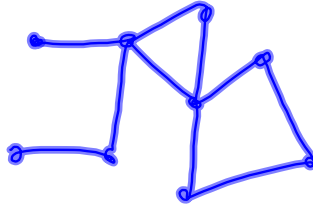


route



lukket rute

Tur = rute uden kant gentagelse



En vej (Path) =  
rute uden pkt gentagelse

$(a,b)$ -vej

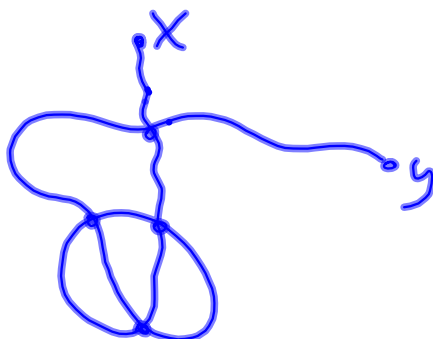


Kreds = 'lukket' vej ( $a=b$ )

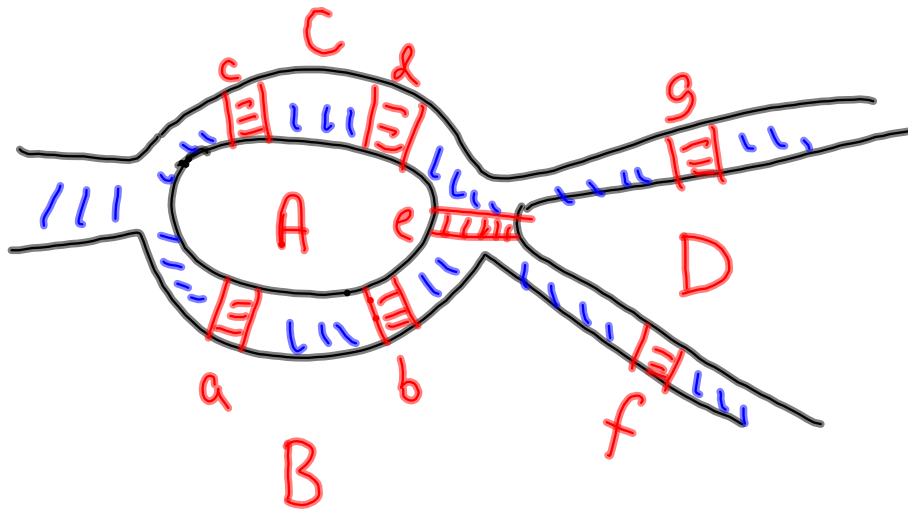
$G=(V,E)$  er sammenhengende  
 $\Updownarrow$  def

$\Updownarrow \forall x,y \in P(G) \exists (x,y)\text{-vej}$

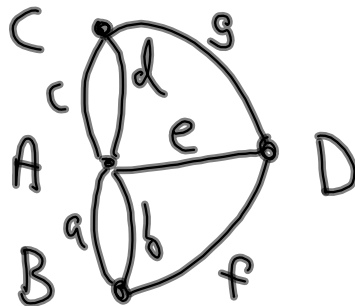
$\Downarrow \forall x,y \in P(G) \exists (x,y)\text{-rute}$



Korteste  $(x,y)$ -rute er en  
 $(x,y)$ -vej



B



Euler-tur = tur som  
indeholder alle kanter i  $G$

Sætning 1.1  $\forall G = (V, E)$  gælder

$$\sum_{x \in V} v(x, \epsilon) = 2|K(G)| = 2|E|$$

B:



## Sætn. 1.2

Enhver graf har et lige antal pkt med ulige valens.

$$\begin{aligned}
 B: \quad 2|E| &= \sum v(x,6) \\
 &= \sum_{\text{ulige Val.}} v(x,6) + \sum_{\text{lige Valens}} v(x,6)
 \end{aligned}$$

Diagram illustrating the proof:

- A blue arrow points from the word "Lige" (Even) to the first summation term  $\sum_{\text{ulige Val.}} v(x,6)$ .
- A green arrow points from the word "Lige" (Even) to the second summation term  $\sum_{\text{lige Valens}} v(x,6)$ .
- A blue arrow points from the word "Lige" (Even) to the entire equation.
- A blue arrow points from the word "Lige" (Even) to the second summation term.



Sætn 1.3 Euler (1736)

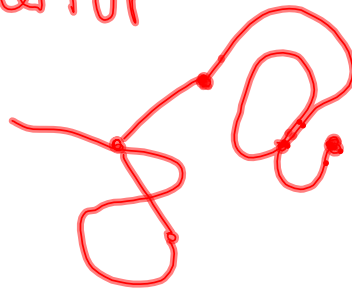
smh  $G$  har en lukket

Euler tur

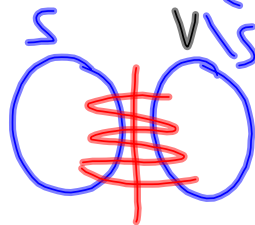


$$\nu(x, G) \equiv 0 \pmod{2} \quad \forall x \in P(G)$$

B:  $\Downarrow$  Lad  $T$  være en lukket  
Eulertur



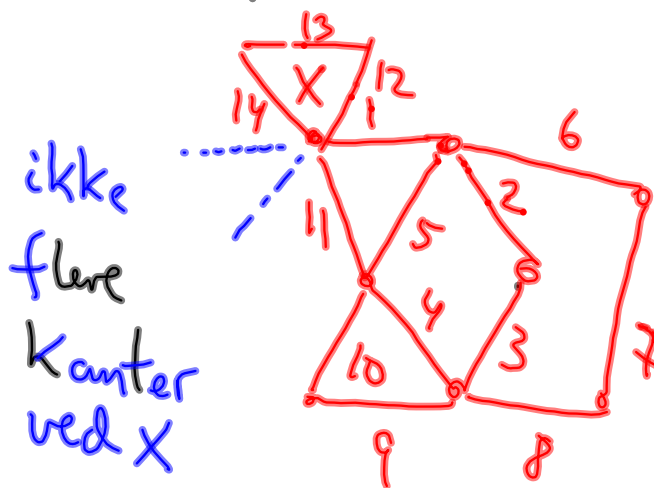
stærkere ( faktisk ækvivalent! )



$T$  har lige antal  
kæmter mellem  
 $S$  og  $VIS$

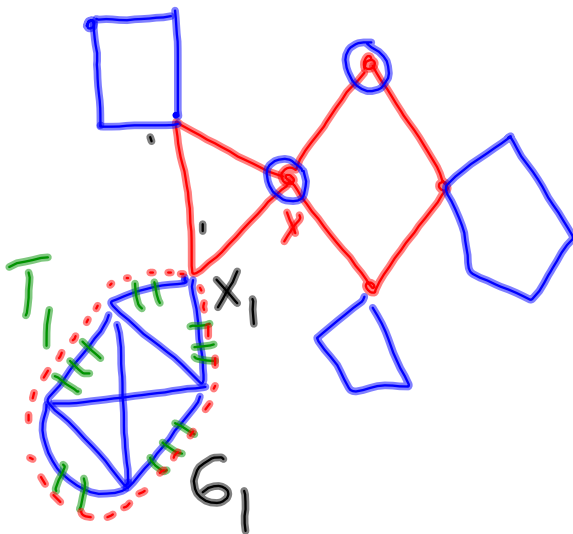
↑↑ Giver et konstruktivt bevis  
(bygger en Euler tur)

- vælg vilk.  $x \in P(G)$
- Lav en maksimal tur  $T$  fra  $x$ :



- Hvis  $T \cong K(G) \checkmark$
- Ellers se på  $G'$

$G' = G$  minus kanterne i  $T$

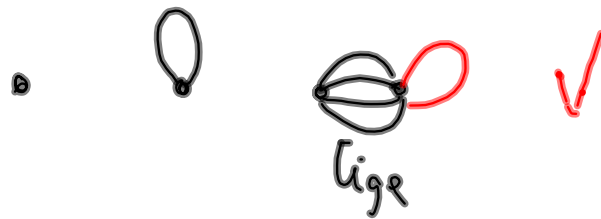


$G'$  har smh  
komponenter  
 $G_1, G_2, \dots, G_k$   $k > 1$   
hvor  $G_k = \{x\}$

alle valenser i  $G_i$   
er lige ( $\forall i$ )

alternativt bevis: induktion  
over  $n = |P(G)|$

basis  $n \leq 2$



ind. antagelse Sætningen  
gælder  $\forall$  smk  $G'$  med  
 $< n$  punkter.  
Se forrige side !!

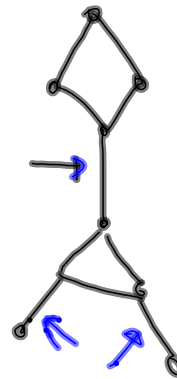
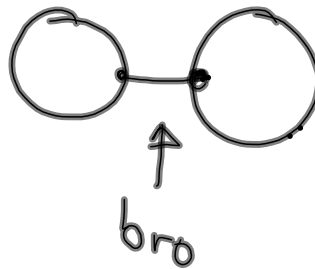
bevis vha maksimalitet  
(modstrid)

antag  $T$  er længst mulig  
(flest mulige kanter)

Vis at  $T$  er Euler  
Se tidligere figur!

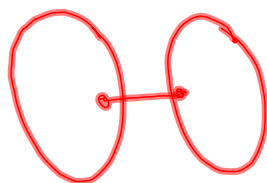
algoritme:

def



observation

Hvis  $v(x, G)$  lige  $\forall x \in P(G)$   
Så har  $G$  ingen broer



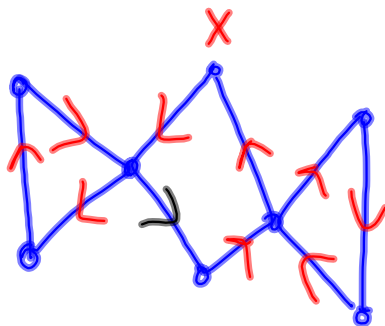
algoritmen vælger nye

Kant fra aktuelle pkt

$x_{i+1}$  ( $x_1 - x_2 - x_3 - \dots - x_i$ )

således, at  $(x_i, x_{i+1})$  ikke  
er en bro i  $G - \{(x_1, x_2), \dots, (x_{i-1}, x_i)\}$

Hvis det er muligt



Antag at  $T$  ikke er

Euler tur  $G' = G - \text{kanter } T$

$$A = \{x \mid v(x, G') = 0\}$$

$$B = \{y \mid v(y, G') > 0\}$$

