

## Cryptology – F03 – Note 9

### Lecture, March 25

We covered the first four sections of chapter 6, introduced digital signature schemes from chapter 7, and motivated the study of hash functions from this.

### Lecture, April 8

We will begin on hash functions from chapter 4, skipping section 4.3.1.

### Lecture, April 15

We will continue with chapter 7.

### Problems for Thursday, April 24

1. Do problem 4.1 in the textbook, but for part (c), the right-hand side of the inequality is wrong. It should be  $2S + N - \frac{N^2}{M}$ . For part (d), use the fact that the left-hand side in (c) is at least zero.
2. Do problem 4.6.
3. Do problem 4.12. For part (b), you can find a (1,1)-forger. Skip the difficult case mentioned.
4. Let  $p$  be an odd prime and  $g_0$  and  $g_1$  be generators of  $\mathbb{Z}_p^*$ . Consider the following two functions:  $f_0(x) = g_0^x \pmod{p}$  and  $f_1(x) = g_1^x \pmod{p}$ . Use these two functions to create a hash function which will hash an arbitrary length message down to a value in  $\mathbb{Z}_p^*$ . Can you make it secure under the assumption that the discrete log problem is infeasible?