# DM508 – Algorithms and Complexity – 2012 Lecture 5

### Lecture, February 10

We continued with NP-Completeness, showing that 3-SAT, CLIQUE, and HAMILTONIAN CIRCUIT are NP-Complete. problems.

## Lecture, February 13

We will cover Cook's Theorem, proving that SATISFIABILITY from the section by Papadimitriou and Steiglitz from the course notes.

### Lecture, February 20

We will finish with NP-Completeness, showing that VERTEX COVER, INDEPENDENT SET, and SUBSET SUM are NP-Complete. We will being on amortized analysis from Chapter 17 of the textbook and Fibonacci heaps from chapter 19.

# Problems to be discussed on February 17

Do problems:

- 1. 34.2-5, 34.2-8, 34.2-10.
- 2. 34.3-2, 34.3-6.
- 3. The following argument is incorrect. Find the most important error. Consider the following algorithm:

```
Input: n \in \mathbb{N}

for i = 2 to n - 1 do

check if i divides n

if it does then output i

endfor

output -1 if no output yet
```

Checking if i divides n can be done in time  $O(\log n)$  via binary search for an integer k such that  $n = i \cdot k$ .

Thus, the total running time is  $O(n \cdot \log n)$  in the worst case. Since  $O(n \cdot \log n) \subset O(n^2)$ , and  $n^2$  is a polynomial, this algorithm runs in polynomial time. Thus, we have an efficient algorithm for factoring,  $O(n \cdot \log n)$ , so we can break RSA, a famous cryptosystem which is believed to be secure.