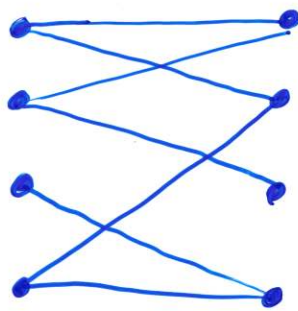


# Bipartite Matching Problem

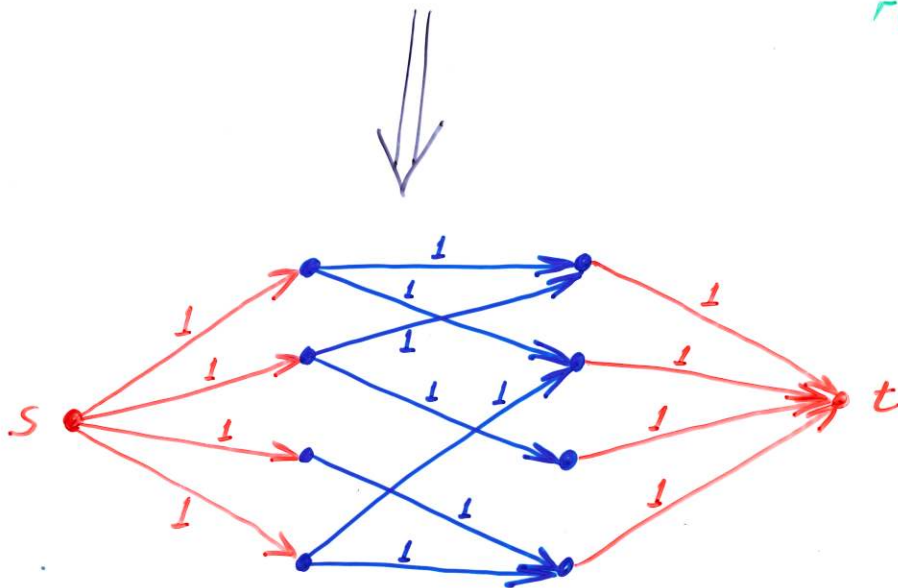
- Cardinality matching

Given a bipartite graph, find as large a matching as possible.  
(No weights, or all weights = 1.)



bipartite matching

reduces to



max-flow  
with capacities  
= 1

Each matching  $\Rightarrow$  flow  
Each integral flow  $\Rightarrow$  matching

$\therefore$  Can compute a maximum matching in time  
 $O(|E|^{3/2})$   
(or even better  $O(|E|\sqrt{|V|})$ .)  
 $O(|V||E|)$  from Ford Fulkerson

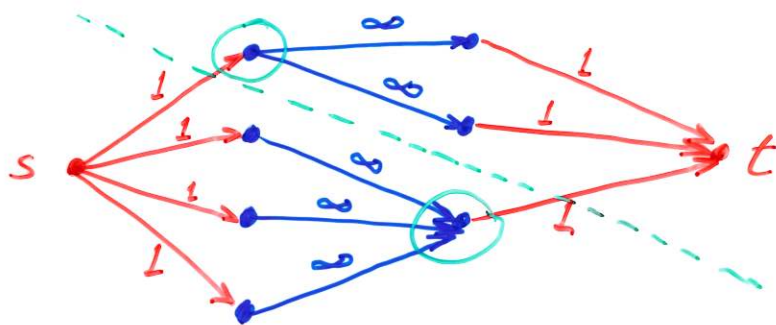
Def  $C = (v_{i_1}, v_{i_2}, \dots, v_{i_k})$  is a cover if every edge is incident with some vertex of  $C$ .

Note For any matching  $M$  and any cover  $C$ ,  
 $|M| \leq |C|$ .

(because, for every edge in  $M$ , any cover must include at least one of the 2 endpoints.)

König-Egervary Thm  $\max_M |M| = \min_C |C|$

Pf



A mincut doesn't include any edges from the original graph. For the cover, choose those vertices ( $\neq s, t$ ) incident to the edge across the cut. Suppose some edge was not covered by this. It would be in the cut. ✗

What about the dual to the linear program for the weighted matching problem?

Suppose the bipartite graph is given by an adjacency matrix.

	$y_1$	$y_2$	$y_3$	$y_4$	$y_5$	$y_6$
$x_1$			X	X	X	
$x_2$		X				
$x_3$	X					
$x_4$		X				
$x_5$		X				
$x_6$		X				

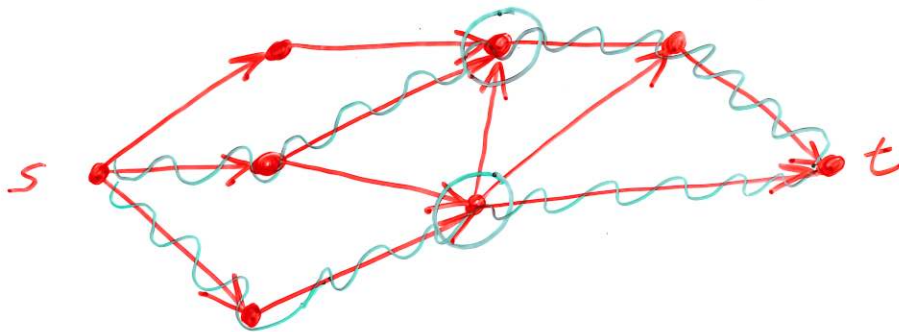
**Matching** - a set of independent  $x$ 's,  
no 2 on the same line (row or column.)

**Cover** - a set of lines that cover all the  $x$ 's,

**König-Egervary Thm** - The maximum # of independent  $x$ 's is = the minimum # of lines to cover all  $x$ 's.

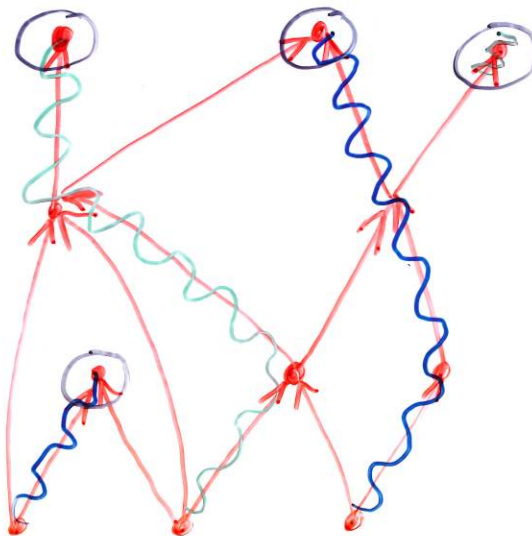
## Other Min-Max Theorems

Menger's Thm If a digraph  $G$  is  $k$ -vertex-connected from  $s$  to  $t$  and does not contain an edge from  $s$  to  $t$ , then  $G$  contains  $k$  independent directed paths from  $s$  to  $t$ .



2

Dilworth's Thm For a given partial order  $(S, \leq)$ , a chain of elements is a sequence  $s_1, s_2, \dots, s_m$  where  $s_i \leq s_{i+1}$ . The minimum number of chains, st. every element is contained in at least one chain, is equal to the maximum number of incomparable elements.



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