

# On-Line Algorithms – F17 – Lecture 1

## Textbook

Allan Borodin and Ran El-Yaniv, *Online Computation and Competitive Analysis*, Cambridge University Press, 1998. There will be supplementary articles.

## Format

The course will be taught by Joan Boyar. The lectures will be in English. The course will generally (there will be a couple of exceptions to these times and places) be at 12:15 on Tuesdays and Wednesdays, and 8:15 on Fridays, in the seminar room. The first two classes will be lectures, but after that we will alternate between lectures and discussion sections. (Ignore what the official schedule about which are lectures and which are discussion sections.) Some classes will be cancelled, since this leads to more than listed in the course description. Be alert for announcements for cancellation. There will be no classes the week starting March 20, but there will be two PhD defenses concerning online algorithms, one on March 21 and the other on March 22, which you are welcome to attend.

There will be assignments which must be approved in order to take the oral exam in June. The assignments are considered “exam projects”. Thus, you may not work with anyone not in your group, but you may work in groups of two to three students. The assignments must be turned in on time. There will be a chance to redo at most two assignments (of the three or four), if they are either late or not good enough the first time. The assignments must be turned in through Blackboard, as one PDF file.

Make sure the names of everyone in your group appears on the cover page, or at the top of the first page. Turn in only one copy per group.

There will be an oral exam in June.

The weekly notes and other information about the course is available through Blackboard or the URL:

<http://www.imada.sdu.dk/~joan/online/>

Please do not hesitate to contact me if you have questions concerning the course. I have office hours on Mondays from 13:00 to 13:45 and on Wednesdays from 9:00 to 9:45.

## Lecture, February 1

We will begin with an introduction to the course. Then, we will begin on chapter 1 in the textbook, probably covering up through section 1.4.

## Lecture, February 3

We will continue with chapter 1 in the textbook.

## Problems for February 7

Note that we may not finish all of these problems. If not, we will continue on February 12.

1. Do Exercise 1.2 in the textbook (assume the static problem).
2. For the proof of Theorem 1.1, work out the details for insertion of an item.
3. This is a modification of Exercise 1.4. The algorithm *fractional MTF* uses a parameter  $d \geq 1$ . Let  $\text{MTF}_d$  be the variant of MTF that behaves as follows. When an item is accessed or inserted at position  $i$ , it moves it  $\frac{i}{d} - 1$  positions closer to the front of the list. Show that  $\text{MTF}_d(\sigma) \leq d(2 \cdot \text{OPT}_C + \text{OPT}_A - n)$ , where  $\text{OPT}_A$  is the total number of transpositions performed by OPT.

Hint: Use  $d$  times the number of inversions for your potential function.

4. Do Exercise 1.6 in the textbook.
5. Consider the algorithm MTF2 for list accessing, which moves an item to the front of the list every second time it accesses it. Let the list consist of the items  $\{a_1, a_2, \dots, a_n, b_1, b_2, \dots, b_n\}$ . Show that on the following

sequence (repeated  $s$  times), the ratio of MTF2's performance to OPT's is at least  $\frac{7}{3}$ .

$$\begin{aligned} &\langle a_1, a_2, \dots, a_n, \quad b_n, b_n, b_{n-1}, b_{n-1}, \dots, b_1, b_1, \\ &\quad a_n, a_{n-1}, \dots, a_1, b_n, b_n, b_{n-1}, b_{n-1}, \dots, b_1, b_1, \\ &\quad b_1, b_2, \dots, b_n, \quad a_n, a_n, a_{n-1}, a_{n-1}, \dots, a_1, a_1, \\ &\quad b_n, b_{n-1}, \dots, b_1, a_n, a_n, a_{n-1}, a_{n-1}, \dots, a_1, a_1 \rangle \end{aligned}$$

Hint: Assume that MTF2 starts with the list ordered as

$$\langle a_1, a_2, \dots, a_n, b_1, b_2, \dots, b_n \rangle$$

while the algorithm it is competing against starts with

$$\langle a_1, a_2, \dots, a_n, b_n, b_{n-1}, \dots, b_1 \rangle$$

Note that this observation is due to Ansgar Grüne and his students and reveals an error in Exercise 1.5 in the textbook.

6. Show that MTF2 is 3-competitive using the potential function method. This can be done counting inversions, as for MTF, but giving the inversions weights based on how many requests to the item are necessary before MTF2 moves it to the front (the weights are 1 or 2). Then the two cases of whether or not MTF2 moves an item to the front are handled separately in the amortized analysis.
7. For the on-line Dual Bin Packing Problem, there are  $n$  bins of unit size. The request sequence contains items of sizes  $0 \leq s \leq 1$  which are to be placed in the bins. The goal is to accept and pack as many items in the request sequence as possible subject to the constraint that no bin should contain items whose sizes add up to more than 1. This is a maximization problem.
  - (a) Give a definition of the competitive ratio for this problem which gives ratios of at most 1.
  - (b) Consider the algorithm First-Fit, which places an item in the first bin in which it fits. First-Fit is a *fair* algorithm, an algorithm which never rejects an item if it fits in one of the bins when it is rejected. Define two more fair algorithms and one which is not fair.

- (c) Assume that  $\frac{1}{k}$  is the size of the smallest item in any sequence. Prove an upper bound of  $\frac{2-\frac{1}{k}}{k}$  on First-Fit's competitive ratio on sequences of this type (assume that OPT must also be fair). Is First-Fit competitive? Prove a lower bound of  $\frac{1}{k}$  on First-Fit's competitive ratio on sequences of this type.