

On-Line Algorithms – F19 – Lecture 12

Lecture, April 1

(On March 27, we reviewed the answers to the problems on the first assignment and did a midway course evaluation.) Continuing with the paper on the relative worst order ratio, we covered Section 3.

Lecture, April 3

We will cover the definitions for relatedness and weakly comparable in Section 2 and Theorem 7 of Section 5. We will begin on the k -server problem from sections 10.1, 10.2, and 10.4 of Chapter 10, in the textbook. We may begin on Chapter 12 in the textbook.

Lecture, April 9

We will finish Section 10.4 and begin on Chapter 12, in the textbook.

Problems for April 8

1. The problems not finished on April 2.
2. What is the complexity of the dynamic programming procedure used for computing the cost of an optimal offline algorithm for the k -server problem when the request sequence is of length n . For the special case of a uniform metric space a faster algorithm exists. What is its complexity?
3. Define and analyze a lazy version of DC for paging.
4. Exercise 10.1.

5. Define a lazy version of DC on the line.
 - (a) Using relative worst order analysis, show that your algorithm, Lazy Double Coverage, is at least as good as Double Coverage (DC) for three points on the line.
 - (b) Show that the greedy algorithm, Greedy, is optimal for three points on the line, according to bijective analysis. To do this, for every algorithm, Alg, you need to find a bijection, f , from sequences of length n to sequences of length n , such that for large enough n , for any sequence I of length n , $\text{Greedy}(I) \leq \text{Alg}(f(I))$.
6. (Easy if you know enough about NP-Completeness) Show that the makespan problem for identical machines is NP-hard.