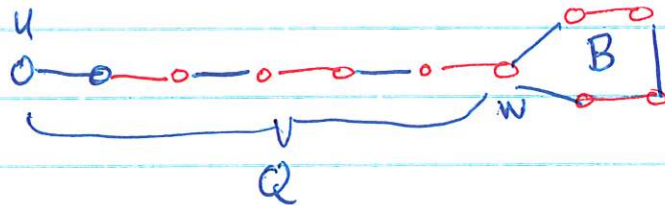
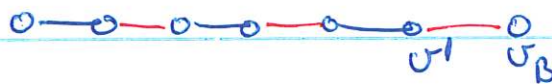


Theorem 10.4 in PS

$G=(V,E)$, M matching, B blossom found while searching from u .



G/B is obtained by contracting B to one vertex v_B :



M/B is M restricted to G/B

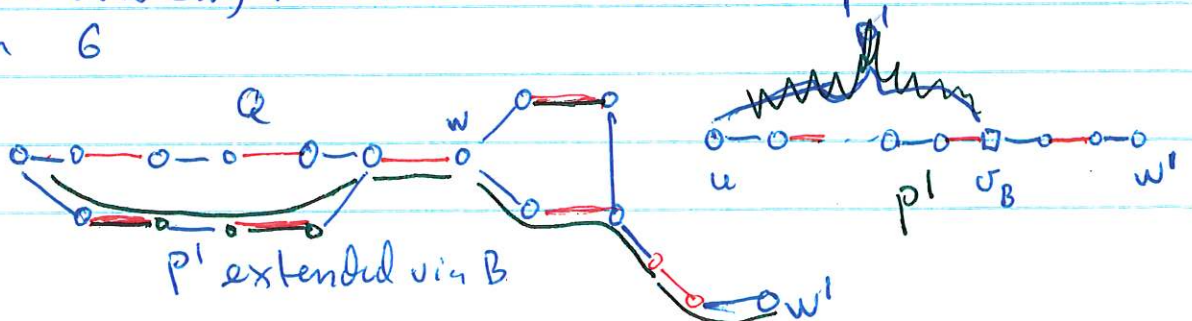
Then the following holds:

G has an M -augmenting path starting at u



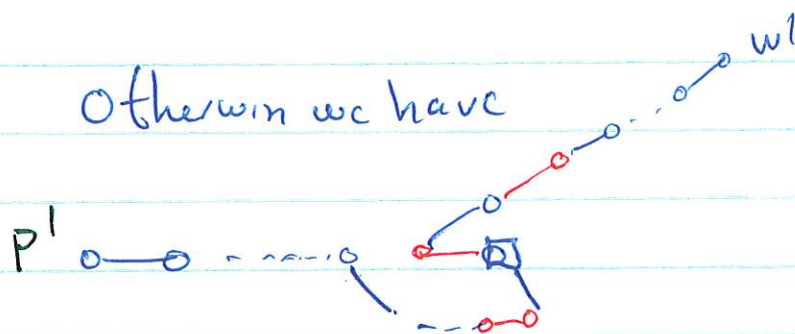
G/B has an M/B -augmenting path starting at u

Proof: $\Uparrow$ let P' be an M/B -augment path in G/B starting at u
 if P' does not contain v_B it is M -augment so
 assume $v_B \in P'$ if P' uses the edge uv_B (see above)
 then it is easy to extend P' to an M -augment path starting at u
 in G

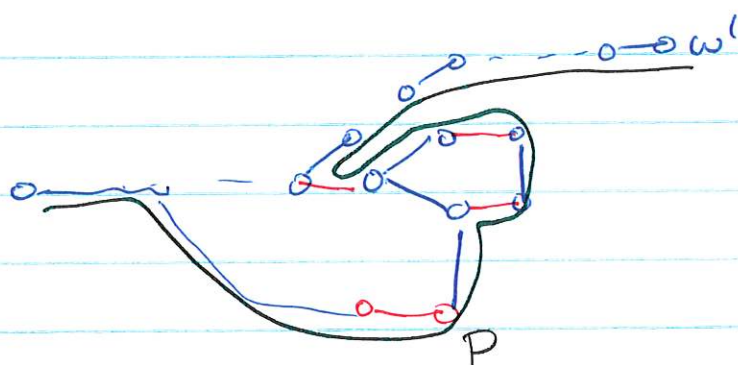


②

otherwise we have



and in G we get the path P :



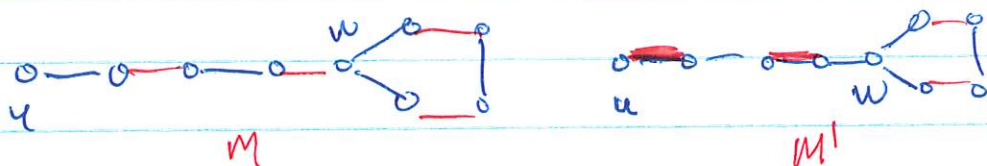
$\Downarrow \exists M$ -augment path starting in u in G

$\Downarrow \exists M^*$ -augment path starting in u in $G^* = P \oplus Q \oplus B$

$\Downarrow$ where M^* is M restricted to G^*

M^* not maximum in G^*

$\Downarrow M'$ not maximum in $G^* \quad M' = M \Delta Q$



$\Downarrow \exists M'$ -augment path in G^* call it P'

$\Downarrow \exists M'/B$ -augment path in G^*/B call it P''
(follow P' from ~~the~~ ~~the~~ the endpoint $u \neq w$
until it hits B for the first time)

③



M'/B not maximum in G^*/B



M/B not maximum in G^*/B

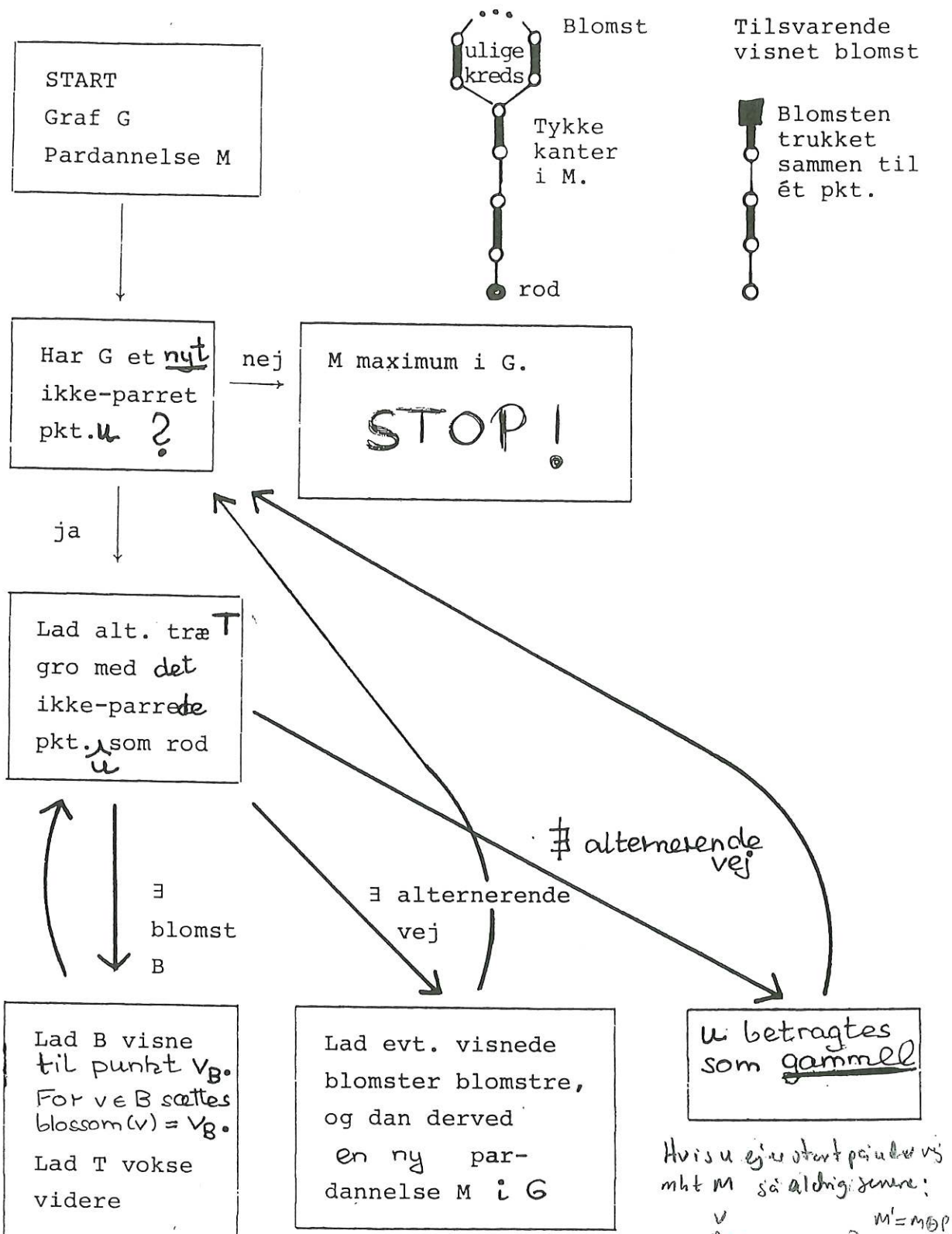


$\exists M/B$ -augmenting path P_B in G^*/B

but G^* has only two unmatched vertices
 u and v and the same holds for
 G^*/B so u is an end vertex of P_B $\square$

a la P&S

Edmond's algoritme for ikke-vægtet generel pardannelse (1965).

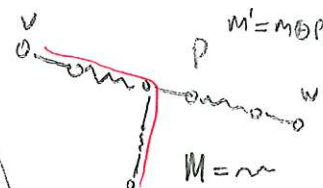


$O(V^3)$ for u enten parret eller gammelt:

max $O(V)$ blomster findes undervejs

$O(V^2)$ arbejde trække blomst sammen hvis bliver en blomst op

selve søgning og udvidelse tager højst $O(V^2)$ for u



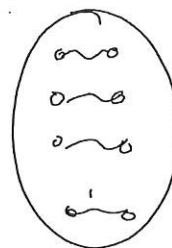
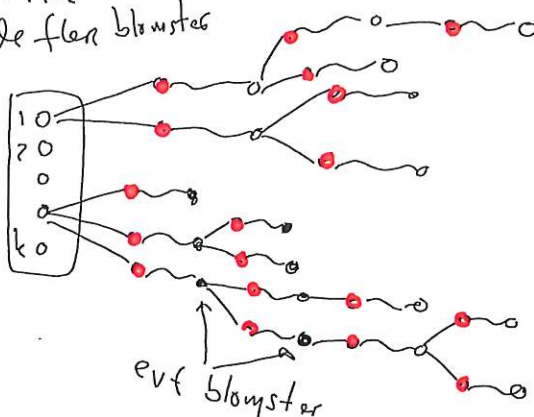
Hvis $\exists$ udv. vej fra u mht $M' = M \oplus P$ så $\exists$ udv. vej fra u mht original M

Sch Thms. i Tutte - Berge

 $\forall$ graph $G=(V, E)$ we have

$$\nu(G) = \min_{u \subseteq V} \frac{1}{2} (|V| + |u| - \text{odd}(G-u))$$

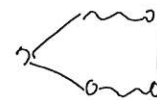
Tuttes sætning ud fra Edmonds alq.

kan ikke
finde flere blomster

U

kan ikke nås.

- alt rd ulige størrelser
- kun forbundet til 1
- da ingen blomster



$$S = \bullet$$

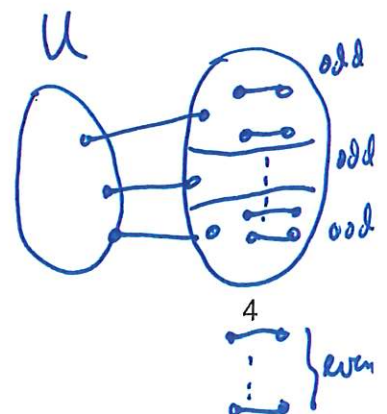
- ulige komponenter i $G-S$

$$\text{odd}(G-S) = |S| + k$$

$$\nu(G) = |M| = \frac{n-k}{2} = \frac{1}{2} (n + |S| - \text{odd}(G-S))$$

Clear that $\nu(G) \leq \min \frac{1}{2} (|V| + |u| - \text{odd}(G-u))$:

$$\begin{aligned} \nu(G) &\leq |u| + \nu(G-u) \\ &\leq |u| + \frac{1}{2} (|V-u| - \text{odd}(G-u)) \\ &= \frac{1}{2} (|V| + |u| - \text{odd}(G-u)) \end{aligned}$$



f-factor

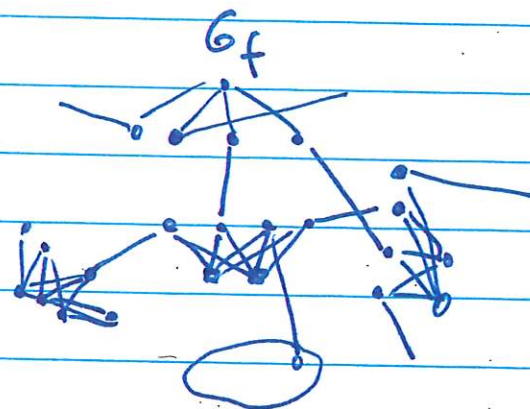
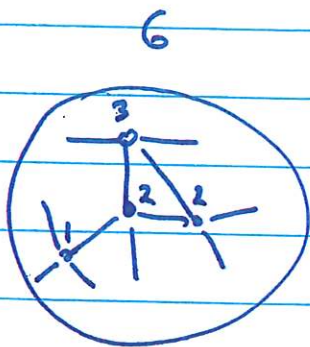
D

$$G = (V, E) \quad f: V \rightarrow \mathbb{Z}_+ \quad f(v) \leq d(v) \quad \forall v$$

Q: does $\exists H = (V, E')$ s.t. $d_H(v) = f(v) \quad E' \subseteq E$

$$|A(v)| = d(v)$$

$$|B(v)| = d(v) - f(v)$$



G_f has 1-factor
 $\Downarrow$
 G has f-factor

Reduction $G \rightarrow G_f$ is polynomial

$$|V(G_f)| = \sum_{v \in V} (2d(v) - f(v)) \leq 2 \sum_{v \in V} d(v) = 4|E|$$

$$|E(G_f)| = |E| + \sum_{v \in V} d(v)(d(v) - f(v))$$

$$\leq |E| + \sum_{v \in V} d(v) \sum_{v \in V} d(v)$$

$$\leq 4|E|^2 + |E|$$