

Theorem 2.1. It is NP-complete to decide whether a 2-regular digraph D contains a pair of arc-disjoint hamiltonian cycles.

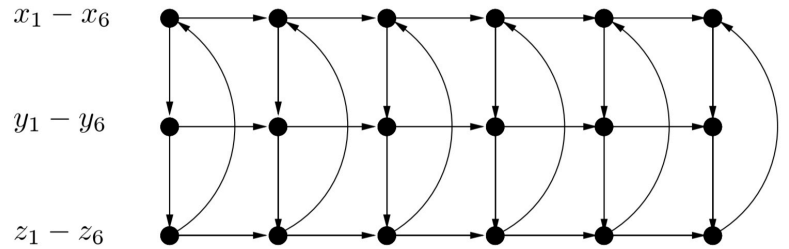
We will reduce the Not-All-Equal 3-SAT problem to this problem.

$$I \text{ 'yes'} \iff D_I \text{ 'yes'}$$

$$H(x, y, z)$$

$$V(H(x, y, z)) = \{x_i, y_i, z_i : i = 1, 2, 3, 4, 5, 6\},$$

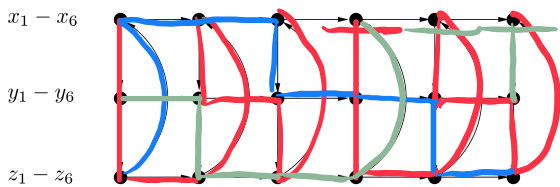
$$A(H(x, y, z)) = \{x_i y_i, y_i z_i, z_i x_i : i = 1, 2, 3, 4, 5, 6\} \cup \{x_j x_{j+1}, y_j y_{j+1}, z_j z_{j+1} : j = 1, 2, 3, 4, 5\}$$



- H-path P start at $x_1 \Rightarrow$ P end at x_6 .

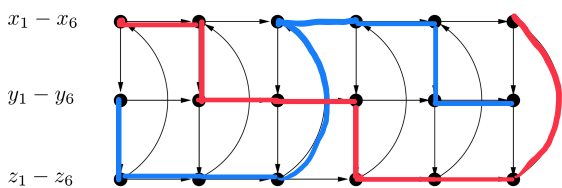
$$Q: y_1 \rightarrow y_6$$

$H(x, y, z) - A(P)$ has a unique 2-path factor QUR. $R: z_1 \rightarrow z_6$



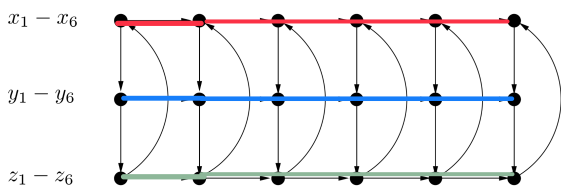
- 2-path factor PUQ start at x_1, y_1 , end in $(x_6, y_6, z_6) \Rightarrow$ P, Q end at x_6, y_6 .

$H(x, y, z) - A(PUQ)$ is a H-path $z_1 \rightarrow z_6$



- 3-path factor PUQUR start at $x_1, y_1, z_1 \Rightarrow$ end at x_6, y_6, z_6 .

$H(x, y, z) - A(PUQUR)$ 6 vertex-disjoint 3-cycles, no arcs between them.



(k-path factor: spanning subdigraph consist of k disjoint path)

Let \mathcal{L} with variables v_1, \dots, v_k and clauses C_1, \dots, C_p .

Assume \mathcal{L} contains v_i and \bar{v}_i , $i \in [k]$.

Construct $D_{\mathcal{L}}$:

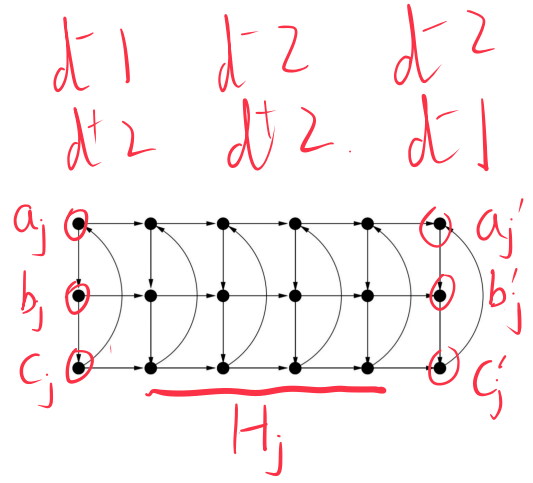
1° \forall clause C_j , let $H_j = H(a_j, a'_j, b_j, b'_j, c_j, c'_j)$

2° \forall variable v , $C_1, \dots, C_p \Rightarrow C_{v,1}, \dots, C_{v,p_v}$

∇ $H_{v,1}, \dots, H_{v,p_v}$
 $H_{\bar{v},1}, \dots, H_{\bar{v},q_{\bar{v}}}$

$1 \leq r \leq p_v - 1$, add an arc from $H_{v,r}$ to $H_{v,r+1}$

$1 \leq r \leq q_{\bar{v}} - 1$, $H_{\bar{v},r}$ $H_{\bar{v},r+1}$



for example:
 $v_1 \begin{cases} H_{v_1,1} := H_1 \\ H_{v_1,2} := H_2 \\ H_{\bar{v}_1,1} := H_3 \end{cases}$

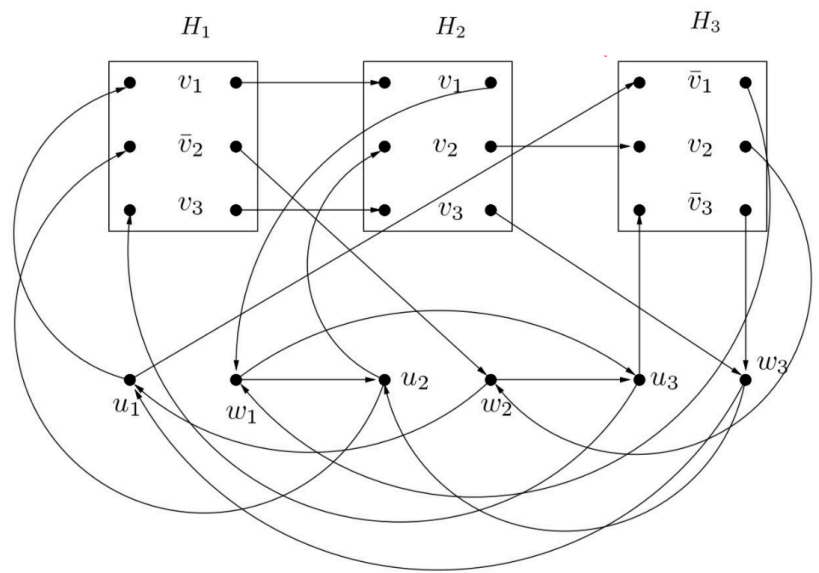
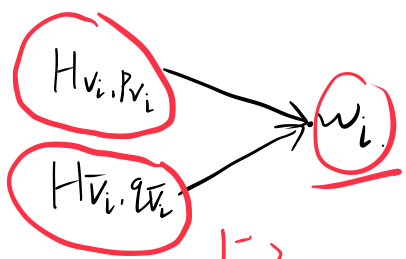
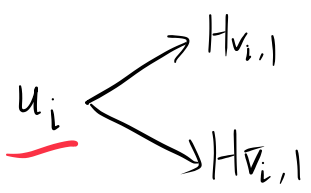


Fig. 4. An illustration of the digraph $D_{\mathcal{L}}$ for the formula $\mathcal{L} = (v_1 \vee \bar{v}_2 \vee v_3)(v_1 \vee v_2 \vee v_3)(\bar{v}_1 \vee v_2 \vee \bar{v}_3)$. For convenience only the six important vertices of each H_i is shown and in the middle column of each H_i we show the 3 literals in the order they appear in C_i . Thus the top literal corresponds to the two top vertices etc.

3° \forall variable v_i , add u_i, w_i .

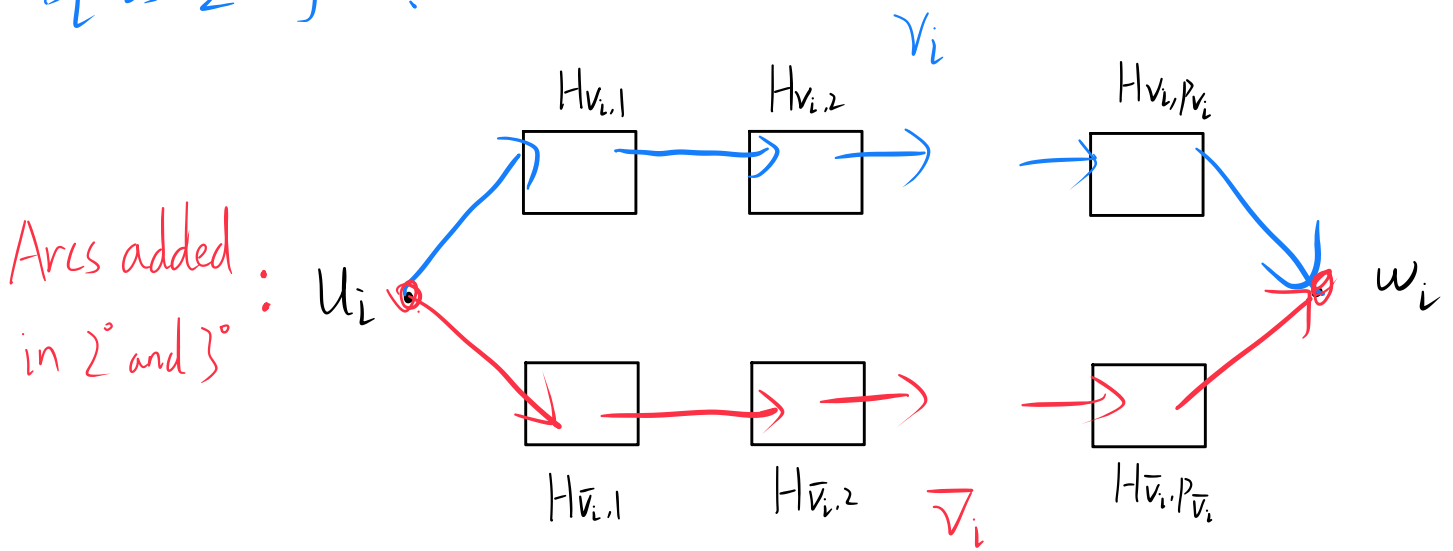


4° $\forall i \in [k]$, add $(w_i, u_{i-1}), (w_i, u_{i+1})$

$d^+ 2$
 $d^- \emptyset 2$

$d^- 2$
 $d^+ \emptyset 2$

D_L is 2-regular.



Claim: L has a satisfying truth assignment t if and only if D_L has a pair of arc-disjoint hamiltonian cycle C, C' .

\Rightarrow :

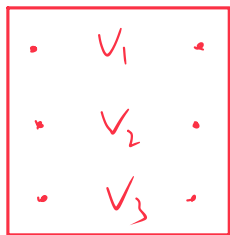
For each variable v_i

v_i is true: let C contain the arcs corresponding to v_i
 C'

v_i is false: let C contain the arcs corresponding to \bar{v}_i
 C'

Then let C contain the arc (w_i, u_{i+1})
 C' (w_i, u_{i-1})

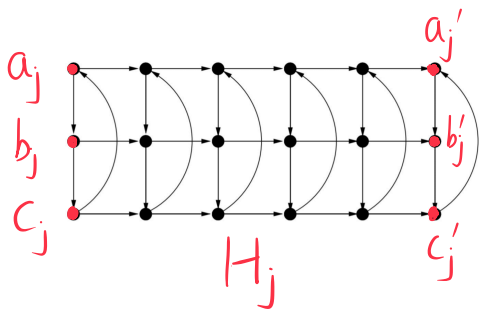
In fact, C choose the path corresponding to true literal,
 C' false



C_j has s true literal, $3-s$ false literal, $1 \leq s \leq 2$
 C choose s -path factor, C' choose $(3-s)$ -path factor

C, C' are closed trail $\xrightarrow{\text{Properties of } H(x,y,z)}$ C, C' are arc-disjoint H -cycles

$\Leftarrow \exists D_j$ has a pair of arc-disjoint hamiltonian cycles C, C' .



C and C' through H_j at least once,
at most twice

If C use the arc from u_i to $H_{v_{i+1}}$, set v_i true.

u_i to $H_{\bar{v}_{i+1}}$, set v_i false.

A literal is true $\iff C$ uses the arcs of D_j that correspond to this literal

So we get truth assignment t .