

DM206 – Advanced Data Structures

Addition to Work Note 3

Defining Asymptotic Notation

Let \mathbb{N} = denote the natural numbers $\{0, 1, 2, \dots\}$ and let \mathbb{R}^+ the positive real numbers.

$$O(f) = \{g: \mathbb{N} \rightarrow \mathbb{R}^+ \mid \exists c \in \mathbb{R}^+ \exists n_0 \in \mathbb{N} \forall n \in \mathbb{N}: n \geq n_0 \Rightarrow g(n) \leq cf(n)\}$$

$$\Omega(f) = \{g: \mathbb{N} \rightarrow \mathbb{R}^+ \mid \exists c \in \mathbb{R}^+ \exists n_0 \in \mathbb{N} \forall n \in \mathbb{N}: n \geq n_0 \Rightarrow g(n) \geq cf(n)\}$$

$$\Theta(f) = O(f) \cap \Omega(f)$$

$$o(f) = O(f) \setminus \Theta(f)$$

$$\omega(f) = \Omega(f) \setminus \Theta(f)$$

$$O(f(m, n)) = \{g: \mathbb{N}^2 \rightarrow \mathbb{R}^+ \mid \exists c \in \mathbb{R}^+ \exists m_0, n_0 \in \mathbb{N} \forall m, n \in \mathbb{N}: m \geq m_0 \wedge n \geq n_0 \Rightarrow g(m, n) \leq cf(m, n)\}$$

[there are many alternative ways of defining asymptotic notation]

Repetition Problems

1. Show that $O(\log_a n) = O(\log_b n)$, where $a, b > 1$.
2. Show that $O(n) \subset O(n \log n) \subset O(n^2)$.
3. Fill in the following table with X's; and arguments.

A	B	$A \in O(B)$	$A \in o(B)$	$A \in \Omega(B)$	$A \in \omega(B)$	$A \in \Theta(B)$
$\log \log n$	$\log n$					
$(\log n)^c$	n^k					
$\frac{\log n}{\log \log n}$	$\log \log n$					
\sqrt{n}	$n^{\sin n}$					
$\log n!$	$\log n^n$					

where c and k are positive constants.

4. Let c, c_1, c_2 be constants. How does $T(n)$ grow asymptotically with the following definitions of T ?
 - (a) $T(n) = T(\frac{n}{2}) + c$
 - (b) $T(n) = 2T(\frac{n}{2}) + c$
 - (c) $T(n) = 3T(\frac{n}{2}) + c$
 - (d) $T(n) = T(\frac{n}{2}) + n$
 - (e) $T(n) = 3T(\frac{n}{2}) + n$
 - (f) $T(n) = T(n - c_1) + c_2$
 - (g) $T(n) = T(n - c) + n$

Assume that n is on some convenient form (a power of two or similar is often helpful) and that $T(1)$ is some (appropriate) constant.