

# Program Extraction from Large Proof Developments

Luís Cruz-Filipe<sup>1,2</sup>

Bas Spitters<sup>1</sup>

<sup>1</sup> University of Nijmegen, Netherlands

<sup>2</sup> Center for Logic and Computation, IST (Lisbon), Portugal

# Why?

- Large constructive library
- Coq has extraction mechanism
- It doesn't work

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# Extraction

- BHK-interpretation: connectives
- Kleene's realizability: a more formal approach
- Curry–Howard isomorphism: proofs  $\longleftrightarrow$  programs
- In practice: algorithm vs. properties; types as “markers”

# FTA-logic

- No elimination of **Prop** terms over **Set**  $\leadsto$  no function definition by cases
- All logic in **Set**
- Extracted program too big

# A solution?

Identify *computationally meaningful* propositions; put everything else in **Prop**.

~> most proof terms can be put back in **Prop**

~> significant amount of “dead code” is eliminated

(for more details see paper in Procs. TPHOLS)

$$\neg : s \rightarrow \mathbf{Prop}$$

$$\rightarrow : s_1 \rightarrow s_2 \rightarrow s_2$$

$$\vee : s_1 \rightarrow s_2 \rightarrow \mathbf{Set}$$

$$\wedge : s_1 \rightarrow s_2 \rightarrow \begin{cases} \mathbf{Prop} & s_1 = s_2 = \mathbf{Prop} \\ \mathbf{Set} & \text{otherwise} \end{cases}$$

$$\forall : \Pi(A : s_1).(A \rightarrow s_2) \rightarrow s_2$$

$$\exists : \Pi(A : \mathbf{Set}).(A \rightarrow s) \rightarrow \mathbf{Set}$$

# Results

- FTA: extracts, compiles, runs... but does not terminate
- Rational numbers: everything is (almost) instantaneous
- Somewhere in between:  $e$ ,  $\pi$  and  $\sqrt{2}$

# Computing $e$

$$e \stackrel{\text{def}}{=} \sum_{n=0}^{+\infty} \frac{1}{n!}$$

~> each term is a rational (constant sequence)

~> but much is going on...

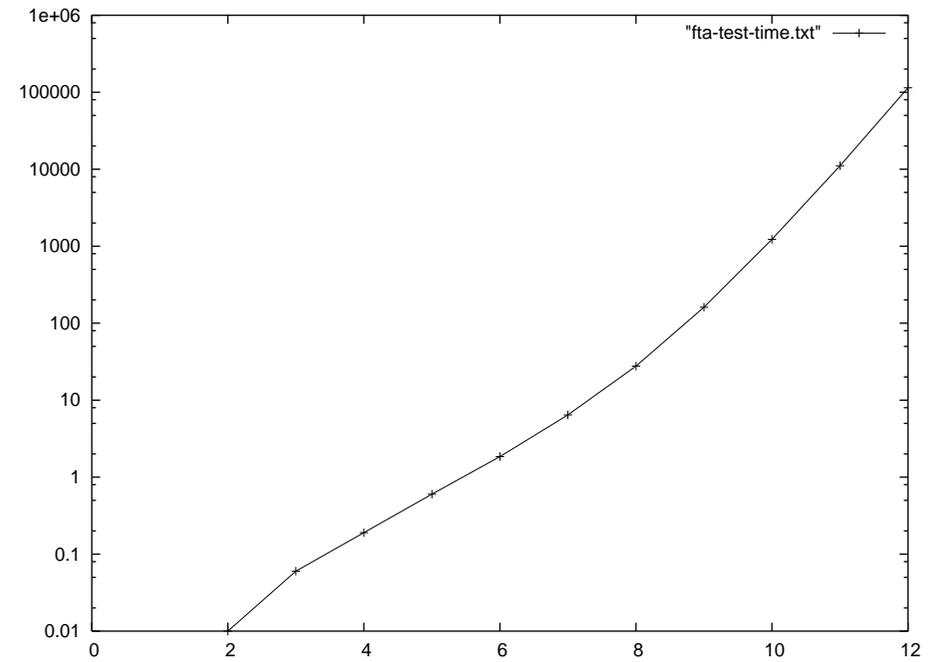
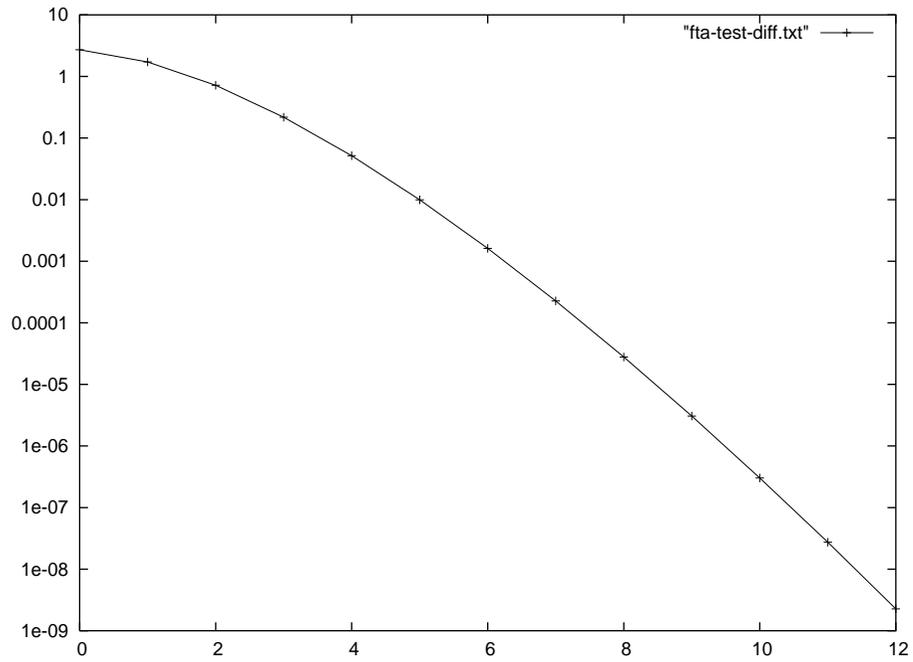
# Immediate Problems. . .

- Unary natural numbers
- A direct proof of  $k! \neq 0$  requires computing  $k!$  in unary notation

## ... & Solutions

- Directly inject  $\mathbb{Z}^+$  into  $\mathbb{R}$
- Prove  $k! \neq 0$  by induction on  $k$

# Some statistics. . .



# Still better

Optimize performance by working directly in the model:

- More efficient definition of factorial
- Simpler proofs and smaller proof terms

~→ 100<sup>th</sup> approximation in 77 seconds (with 157 correct digits)

# Conclusions

- The more abstract the formalization, the less efficient the extracted program
- Obtaining a *working* program is far from straightforward
- Small, carefully thought, modifications in the formalization can make huge differences in the extracted program
- Future improvements in Coq may also make huge differences. . .

# Future Work

- A similar analysis of  $\sqrt{2}$
- Improving the extraction mechanism: pruning, modules
- Eventually: the FTA