

active integrity constraints for multi-context systems

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structure

- 1 motivation
- 2 the context
- 3 our formalism
- 4 evaluation

integrity constraints in reasoning systems

relational dbs

the classical setting

deductive dbs

mostly mid-1980s

- ↪ separate integrity constraints from data
- ↪ integrity constraints as preferred models

ontologies

last 15–20 years

- ↪ open-world semantics makes the problem different
- ↪ integrity constraints as terminological axioms
(but with a different semantics)

*heterogeneous
systems*

last 10 years

- ↪ in multi-context systems (our setting)
- ↪ internalize integrity constraints

in general

no continuation, no apparent consensus

our goal

- generalize existing notions in particular frameworks (e.g. relational databases)
- expressive enough to capture conditions spanning several systems
- decidability? good complexity bounds?
- algorithms for repairing inconsistencies

our target

- active integrity constraints (flesca *et al.*, '04)
- defined for relational databases
- allow to express both *constraints* and *repair actions*
- “good” algorithms for repair
- ... and it's kind of a nice formalism

our contribution

- active integrity constraints in a general-purpose framework
- captures previous constructions as special cases
- clean separation between consistency and integrity
- including repair actions avoids need for abduction
- repairs can be computed automatically (with a grain of salt)

active integrity constraints (for databases)

main idea

datalog-style rules

- body specifies an integrity constraint (clausal, denial form)
- heads are sets of “repair actions” (alternative)
- several different semantics

algorithms

tree-based algorithms

- can compute different kinds of repairs
- (non-deterministic) polynomial complexity
- sometimes require extra testing (complexity...)

multi-context systems

main idea

- reasoning systems (“contexts”)
- connected by datalog-style rules (“bridge rules”)

brewka & eiter
'07

heterogenous non-monotonic multi-context systems

heterogeneous contexts can use different logics

non-monotonic bridge rules can contain negation

multi-context several different systems

equilibrium

an equilibrium is a set of beliefs that is compatible with all knowledge bases and bridge rules

↪ think logic programming. . .

ontologies

in particular, we can view an ontology as an mcs

- the a-box is one context
- the t-box is another context
- there are bridge rules injecting all instances from the a-box to the t-box
- equilibria are sets all queries that return true

↪ separation between a-box and t-box is useful for some types of integrity constraints

active integrity constraints

syntax

an active integrity constraint is written as a bridge rule with disjunctive head

validity

actions on the head must satisfy some constraints

- for every action, there is an inconsistent state it repairs
- for every inconsistent state, there is an action that repairs it



validity is undecidable in general, arguably simple in practice

evaluation

- capture all classes of ontology integrity constraints from a 2013 survey
- are able to define actions for their heads and show validity
- examples in paper, but approach is systematic
- also discuss some other types of ontology integrity constraints

some examples

*specific type
constraints*

$(A:\text{gradStudent}(X)), (A:\text{student}(X))$
 $\implies (A:\text{del}(\text{student}(X)))$

*property
domain
constraints*

$(T:\text{enrolled}(X,Y), \text{not } (T:\text{student}(Y)))$
 $\implies (A:\text{add}(\text{student}(Y)))$

*functional
dependencies*

$(A:\text{hasEmail}(X,Z)), (A:\text{hasEmail}(Y,Z)), \text{not } (T:(X=Y))$
 $\implies (A:\text{del}(\text{hasEmail}(X,Z))) \mid (A:\text{assert}(X=Y))$

*minimum
cardinality
constraints*

$(T:(\leq 10.\text{enrolled})(X))$
 $\implies (A:\text{redistribute}(\neg\text{class}(X)))$

thank you!