

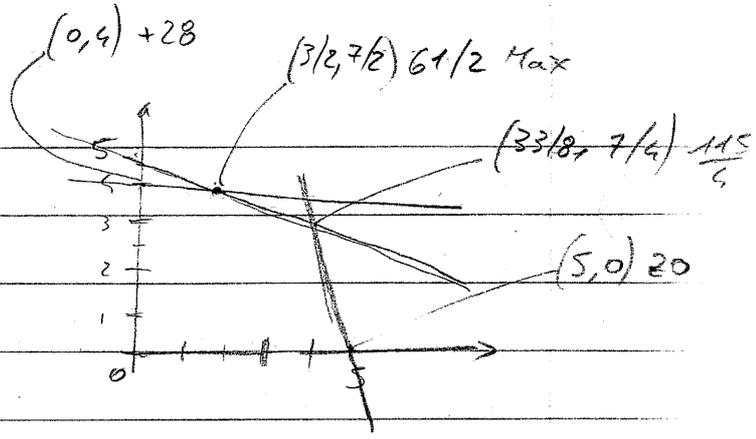
TASK 1

$$\max z = 4x_1 + 7x_2$$

$$x_1 + 3x_2 \leq 12$$

$$-4x_1 + 6x_2 \leq 27$$

$$4x_1 + 2x_2 \leq 20$$



(a) (0,0) is feasible $\Rightarrow z=0=LB$

	x_1	x_2	s_1	s_2	s_3	$-z$	b	b/a_{i1}	b/a_{i2}	$c_j \cdot b/a_{ij}$
(c)	1	3	1	0	0	0	12	12	4	$7 \cdot 4 = 28$
(d)	4	6	0	1	0	0	27	$27/4 = 8$	$27/6 = 4.5$	
	4	2	0	0	1	0	20	5	10	$5 \cdot 4 = 20$
	4	7	0	0	0	-1	0			

largest coefficient $\rightarrow 7$

largest increase $\rightarrow 7$

$1/3$	1	$1/3$	0	0	0	0	4
2	0	-2	1	0	0	0	3
$10/3$	0	$-2/3$	0	1	0	0	12
$-5/3$	0	$7/3$	0	0	1	0	28

(e)	0	1	$2/3$	$-1/6$	0	0	$7/2$
	1	0	-1	$1/2$	0	0	$3/2$
	0	0	$8/3$	$-5/3$	1	0	7
	0	0	$-2/3$	$-5/6$	0	1	$-6\frac{1}{2}$

$$GAP = \frac{61/2 - 0}{61/2} = 1$$

(f) I row: $\frac{2}{3}S_1 + \frac{5}{6}S_2 \geq \frac{1}{2}$ $\frac{1}{2}S_1 \geq \frac{1}{2}$

$$\frac{2}{3}(12 - x_1 - 3x_2) + \frac{5}{6}(27 - 4x_1 - 6x_2) \geq \frac{1}{2}$$

$$8 - \frac{2}{3}x_1 - 2x_2 + \frac{135}{6} - \frac{10}{3}x_1 - 5x_2 \geq \frac{1}{2}$$

$$-4x_1 - 7x_2 \geq \frac{1}{2} - 8 - \frac{135}{6}$$

II row: $\frac{1}{2}S_2 \geq \frac{1}{2}$ (i)-form

$$\frac{1}{2}(27 - 4x_1 - 6x_2) \geq 1$$

$$-4x_1 - 6x_2 \geq -26$$

$$-2x_1 - 3x_2 \geq -13$$

$$2x_1 + 3x_2 \leq 13$$

(ii)-form

(g)

Introducing the i-form we have a basis but infeasible -

Introducing the ii-form we do elementary row operations to arrive to a canonical form with infeasible basis -

We use dual simplex:

$$\begin{array}{cccc|cccc}
 0 & 0 & 0 & -\frac{1}{2} & 0 & 0 & 1 & -\frac{1}{2} \\
 0 & 1 & \frac{2}{3} & -\frac{1}{6} & 0 & 0 & 0 & \frac{7}{2} \\
 1 & 0 & -1 & \frac{1}{2} & 0 & 0 & 0 & \frac{3}{2} \\
 \hline
 0 & 0 & \frac{8}{3} & -\frac{5}{3} & 1 & 0 & 0 & 7 \\
 0 & 0 & -\frac{2}{3} & -\frac{5}{6} & 0 & 1 & 0 & -\frac{61}{2}
 \end{array}$$

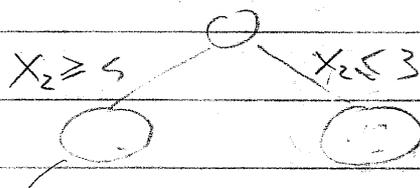
1) pivot < 0

2) row with b term negative

3) col that $\min \left| \frac{c_j}{a_{ij}} \right| \Rightarrow -\frac{1}{2}$ is pivot

(h)

Using branch and bound on we have



we could use heuristics again at each node

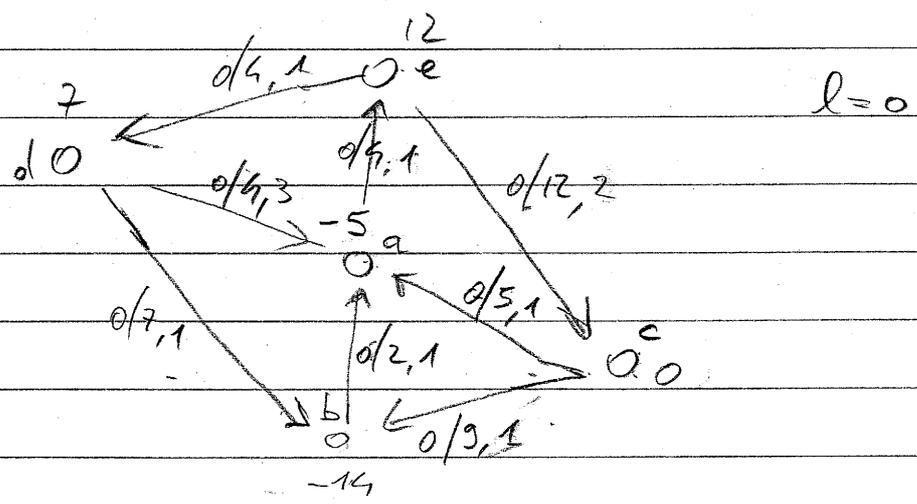
the LP requires a dual-simplex step sol. will be

(0, 4) of val 28

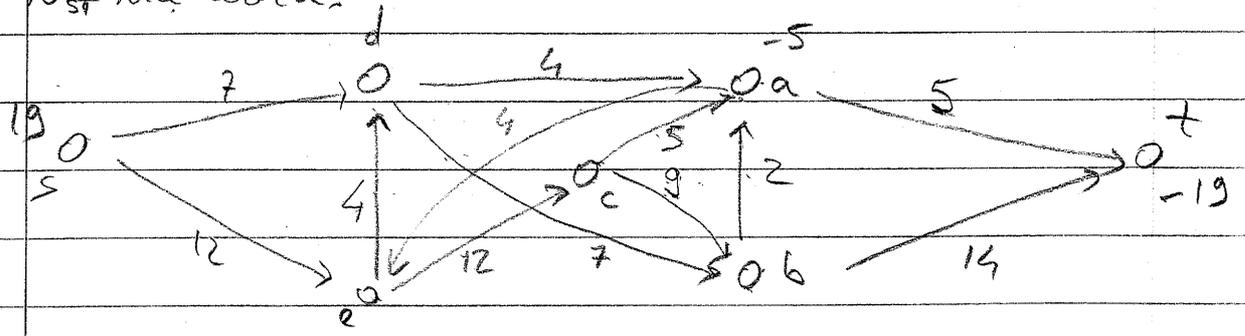
sol will be (2, 3) of val 29

TASK 2

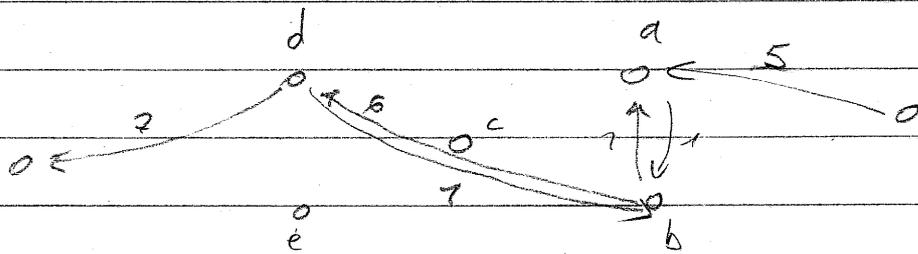
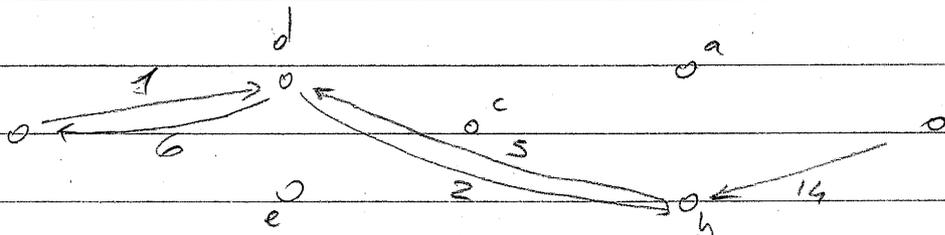
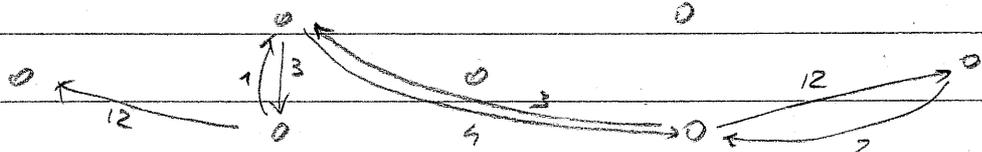
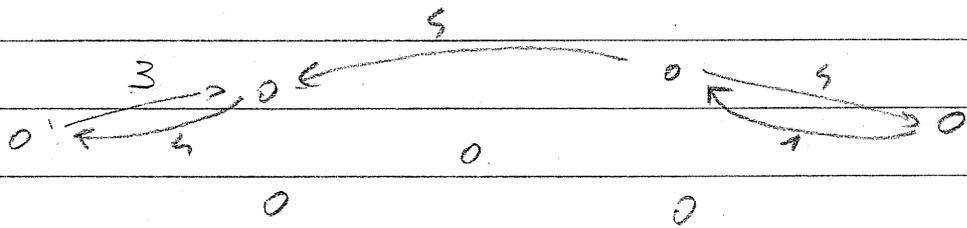
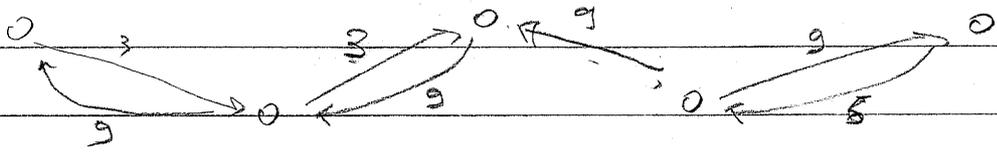
(a)



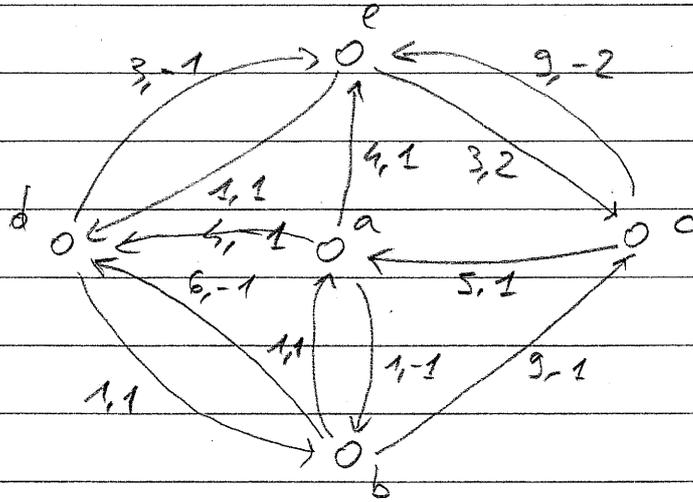
Net network:



Apply max flow to N_{st} network. Ford Fulkerson alg:

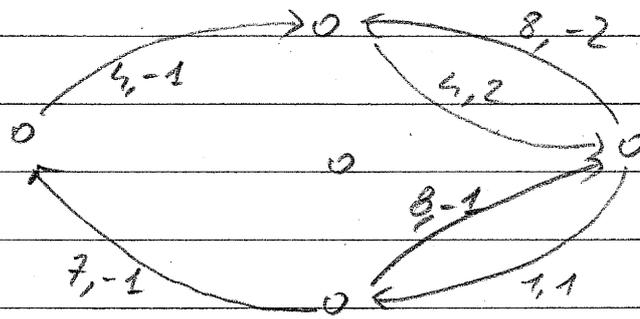


(b) $N(x)$

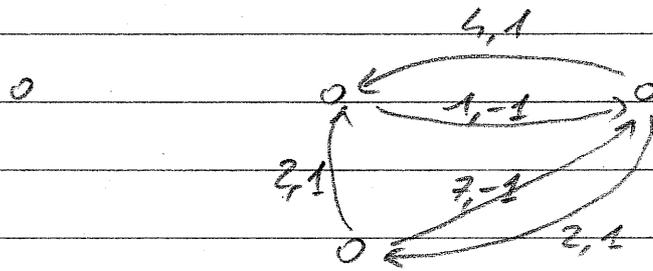


cost 41

(c)

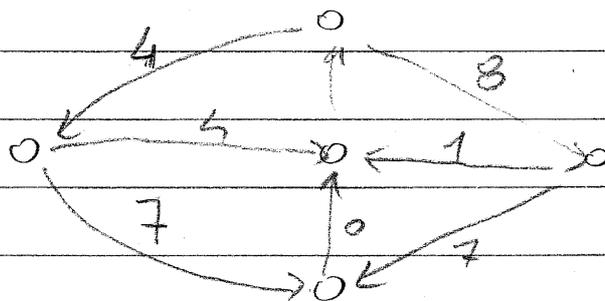


cost 40



cost 39

final flow



TASK 3

(B)

$$\boxed{A_B \quad A_N \quad b}$$

$c_B \quad c_N$

$$A_B X_B + A_N X_N = b$$

$$X_B = A_B^{-1} b$$

$$\bar{c}_B^T = c_B^T - \pi A_B = 0 \Rightarrow c_B^T A_B^{-1} = \pi$$

$$\bar{c}_N^T = c_N^T - \pi A_N = c_N^T - c_B^T A_B^{-1} A_N$$

$$\begin{bmatrix} & x_2 & x_3 & & & & s_3 & & \\ 2 & 1 & 5 & 1 & 1 & 0 & 0 & 0 & 8 \\ 2 & 2 & 0 & 4 & 0 & 1 & 0 & 0 & 12 \\ 3 & 1 & 2 & 0 & 0 & 0 & 1 & 0 & 18 \\ 1 & 2 & 1 & 1 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

$$A_B = \begin{bmatrix} 5 & 1 & 0 \\ 0 & 2 & 0 \\ 2 & 1 & 1 \end{bmatrix}$$

$$A_B^{-1} = \begin{bmatrix} \frac{2}{10} & 0 & -\frac{1}{10} \\ -\frac{1}{10} & \frac{5}{10} & -\frac{3}{10} \\ 0 & 0 & 1 \end{bmatrix}^T = \begin{bmatrix} 0,2 & -0,1 & 0 \\ 0 & 0,5 & 0 \\ -0,4 & -0,3 & 1 \end{bmatrix}$$

$$X_B = \begin{bmatrix} 0,2 & -0,1 & 0 \\ 0 & 0,5 & 0 \\ -0,4 & -0,3 & 1 \end{bmatrix} \begin{bmatrix} 8 \\ 12 \\ 18 \end{bmatrix} = \begin{bmatrix} 1,6 - 1,2 \\ 6 \\ -3,2 - 3,6 + 18 \end{bmatrix} = \begin{bmatrix} 0,4 \\ 6 \\ 11,2 \end{bmatrix} = \begin{bmatrix} x_3 \\ x_2 \\ s_1 \end{bmatrix}$$

$$\bar{c}_N^T = [1 \ 1 \ 0 \ 0] - [1 \ 2 \ 0] \begin{bmatrix} 0,2 & -0,1 & 0 \\ 0 & 0,5 & 0 \\ -0,4 & -0,3 & 1 \end{bmatrix} \begin{bmatrix} 2 & 1 & 1 & 0 \\ 2 & 4 & 0 & 1 \\ 3 & 0 & 0 & 0 \end{bmatrix} =$$

$$= \begin{bmatrix} 0,4 - 0,2 & 0,2 - 0,4 & 0,2 & -0,1 \\ 1 & 2 & 0 & 0,5 \\ -0,8 - 0,6 + 3 & -0,4 - 1,2 & -0,4 & -0,3 \end{bmatrix}$$

$$= [1 \ 1 \ 0 \ 0] - [1 \ 2 \ 0] \begin{bmatrix} 0,2 & -0,2 & 0,2 & -0,1 \\ 1 & 2 & 0 & 0,5 \\ 1,6 & -1,6 & -0,4 & -0,3 \end{bmatrix} =$$

$$= [1 \ 1 \ 0 \ 0] - [0,2+2 \quad -0,2+4 \quad 0,2 \quad -0,1+1] =$$

$$= [-1,2 \quad -2,8 \quad -0,2 \quad -0,9]$$

All red. costs are negative \Rightarrow sol. is optimal.

$$b) \quad \bar{\pi}^T = c_B^T A_B^{-1} = [1 \ 2 \ 0] \begin{bmatrix} 0,2 & -0,1 & 0 \\ 0 & 0,5 & 0 \\ -0,4 & -0,3 & 1 \end{bmatrix} = [0,2 \ 0,9 \ 0]$$

The dual variables are the reduced costs of slack variables with change of sign. (This has been shown in the proof of the Strong duality Theorem)

$$y^T = [0,2 \quad 0,9 \quad 0] = \bar{\pi}^T$$

c) Looking at the dual var., the first and second constraints are binding. The third has a slack:

$$b_i - \sum_{j=1}^n a_{ij} x_j > 0$$

$$d) \quad z^T = y^T b = [0,2 \ 0,9 \ 0] \begin{bmatrix} 8 \\ 26 \\ 18 \end{bmatrix} = 1,6 + \quad + 0 = 23,4$$

Innovol

before it was

$$c^T x^* = [1 \ 2 \ 1 \ 1 \ 0 \ 0] \begin{bmatrix} 0 \\ 6 \\ 0,4 \\ 0 \\ 0 \\ -11,2 \end{bmatrix} = 12 + 0,4 = 12,4$$

hence obj increases

But we must make sure the b term does not become negative.

$$\bar{b} = A_B^{-1} b = \begin{bmatrix} 0,2 & -0,1 & 0 \\ 0 & 0,5 & 0 \\ -0,4 & -0,3 & 1 \end{bmatrix} \begin{bmatrix} 8 \\ 26 \\ 18 \end{bmatrix} = \begin{bmatrix} 1,6 & -2,6 \\ 13 \\ -3,2 & -7,8 + 18 \end{bmatrix} = \begin{bmatrix} 1,6 & -2,6 \\ 13 \\ 14,8 & -0,8 \end{bmatrix}$$

one term becomes negative \Rightarrow we need to iterate.

② we look back at how we computed red costs.

$$\bar{c}_N = [\textcircled{3} \quad 1 \quad 0 \quad 0] - [2,2 \quad 3,8 \quad 0,2 \quad 0,9] = [0,8 \quad \dots]$$

the red cost of x_1 becomes positive $\Rightarrow x_1$ enters the basis.

To determine the leaving var we need to see how much we can increase:

$$X_B = X_B^* - A_B^{-1} A_N X_N = X_B^* - A_B^{-1} a_e X_N$$

$$A_B^{-1} \cdot a_e = \begin{bmatrix} 0,2 & -0,1 & 0 \\ 0 & 0,5 & 0 \\ -0,4 & -0,3 & 1 \end{bmatrix} \begin{bmatrix} 2 \\ 2 \\ 3 \end{bmatrix} = \begin{bmatrix} 0,4 - 0,2 \\ 1 \\ -0,8 - 0,6 + 3 \end{bmatrix} = \begin{bmatrix} 0,2 \\ 1 \\ 1,6 \end{bmatrix}$$

\downarrow entering col

$$X_B = \begin{bmatrix} 0,4 \\ 6 \\ 11,2 \end{bmatrix} - \begin{bmatrix} 0,2 \\ 1 \\ 1,6 \end{bmatrix} t \geq 0 \quad \begin{array}{l} t \leq 0,4/0,2 = 0,2 \leftarrow \\ t \leq 6 \\ t \leq 11,2/1,6 = 7 \end{array}$$

x_3 leaves the basis

8) We need to determine the red. cost for the new var.

$$c_N^T = c_N^T - \pi A_N =$$

$$= [\textcircled{5} \ 1 \ 1 \ 0 \ 0] - [0,2 \ 0,9 \ 0] \cdot \begin{bmatrix} \textcircled{3} & 2 & 1 & 1 & 0 \\ 5 & 2 & 4 & 0 & 1 \\ 6 & 3 & 0 & 0 & 0 \end{bmatrix} =$$

$$= 5 - 0,6 - 4,5 > 0$$

$$\delta > 5,1$$

TASK 4

(a)

$$A = \begin{array}{cccccc|c} & x_1 & x_2 & x_3 & x_4 & x_5 & y & \\ \hline & & & & & & 1 & \leq 1 \\ & & & & & & -1 & \leq 0 \\ & & & & & & -1 & \leq 0 \\ & & & & & & -1 & \leq 0 \\ & & & & & & -1 & \leq 0 \\ & & & & & & 1 & -1 & \leq 0 \end{array}$$

Δ^T is $\begin{bmatrix} 1 & & & & & & & & \\ & 1 & & & & & & & \\ & & 1 & & & & & & \\ & & & 1 & & & & & \\ & & & & 1 & & & & \\ & 1 & -1 & -1 & -1 & -1 & & & \end{bmatrix} \Rightarrow$ TUM for $I_1 = \text{all}$, $I_2 = \emptyset$

$$\sum_{i \in I_1} a_{ij} = \sum_{i \in I_2} a_{ij} = 0 \quad \forall j$$

the first column has only one entry not zero. However since the matrix without that column is TUM, all its square submatrices have determinant 0,1,-1, hence also all matrices obtained by adding the first column will have determinants 0,1,-1 since all its elements are 0 or 1.

$$b) \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & -1 & 0 & 0 & 0 & 1 & 0 \\ 2 & 0 & 1 & 1 & 1 & 0 & -1 & 0 & 0 & 2 & 0 \\ 3 & 1 & 1 & 1 & 1 & 0 & 0 & -1 & 0 & 3 & 0 \\ 4 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 5 \\ 5 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

for each $i=4,3,2$
 subtract i th row with
 $(i+1)$ th row

$$\begin{array}{cccccccc} 2 & 0 & 1 & 1 & -1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & -1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 & 1 & -1 & 0 & 1 & 0 \\ -1 & 0 & 0 & -1 & 0 & 0 & 1 & -1 & -1 & -1 \\ 0 & -1 & -1 & 0 & 0 & 0 & 0 & 1 & -1 & -1 \end{array}$$

\Rightarrow TUM and min cost flow

$$c) \min \quad 720x_1 + 800x_2 + 760x_3 + 680x_4 + 720x_5 + 780x_6 + 640x_7$$

$$0-6: \quad x_1 + x_2 \geq 2$$

$$6-8: \quad x_2 + x_3 \geq 8$$

$$8-11: \quad x_3 + x_4 \geq 5$$

$$11-14: \quad x_3 + x_4 \geq 7$$

$$14-16: \quad x_4 + x_5 \geq 3$$

$$16-18: \quad x_5 + x_6 \geq 4$$

$$18-20: \quad x_5 + x_6 + x_7 \geq 6$$

$$20-22: \quad x_5 + x_6 + x_7 \geq 3$$

$$22-24: \quad x_6 + x_7 \geq 1$$

$$x_1, x_2, \dots, x_7 \geq 0 \text{ and integer.}$$

The matrix has consecutive 1's property on cols
 hence LP relax gives integer results.

this solution contains the mistakes indicated in red.
See next page for correct solution.

TASK 5

(a)

$$\max x_1 - 2x_2$$

$$x_1 + 2x_2 - x_3 \geq 0$$

(P)

$$4x_1 + 3x_2 + 4x_3 \leq 3$$

$$2x_1 - x_2 + 2x_3 = 1$$

$$x_2, x_3 \geq 0$$

$$P(y_1, y_2, y_3) = \max x_1 - 2x_2 + y_1(x_1 + 2x_2 - x_3) + y_2(4x_1 + 3x_2 + 4x_3 - 3) + y_3(2x_1 - x_2 + 2x_3 - 1)$$

$$\forall y_1 \geq 0, y_2 \leq 0, y_3 \in \mathbb{R} \quad P(y_1, y_2, y_3) \geq \text{opt}(P)$$

$$\min_{y_1, y_2, y_3 \in \mathbb{R}} P(y_1, y_2, y_3)$$

$$\begin{aligned} \max & (1 + y_1 + 4y_2 + 2y_3) x_1 \\ & (-2 + 2y_1 + 3y_2 - y_3) x_2 \\ & (-y_1 - 4y_2 + 2y_3) x_3 \\ & -3y_2 - y_3 \end{aligned}$$

+

$$\max c_1 x_1 + c_2 x_2 + \dots \Rightarrow \text{if } c_i > 0 \Rightarrow x_i = +\infty$$

$$c_i < 0 \Rightarrow x_i = 0$$

Hence:

if $x \in \mathbb{R} \Rightarrow c_i = 0$ or else unbounded

$$\min -3y_2 - y_3$$

$$y_1 + 4y_2 + 2y_3 \leq -1$$

$$2y_1 + 3y_2 - y_3 \leq 2$$

$$-y_1 - 4y_2 + 2y_3 = 0$$

$$y_1 \geq 0 \quad y_2 \leq 0 \quad y_3 \geq 0$$

TASK 5

(a)

$$\begin{aligned} \max \quad & x_1 - 2x_2 \\ & x_1 + 2x_2 - x_3 \geq 0 \\ & 4x_1 + 3x_2 + 4x_3 \leq 3 \\ & 2x_1 - x_2 + 2x_3 = 1 \\ & x_2, x_3 \geq 0 \\ & x_1 \in \mathbb{R} \end{aligned}$$

(P)

We want an UB; we bring (relax) constraints in obj. func.

$$P(y_1, y_2, y_3) = \max x_1 - 2x_2 + y_1(x_1 + 2x_2 - x_3) + y_2(4x_1 + 3x_2 + 4x_3 - 3) + y_3(2x_1 - x_2 + 2x_3 - 1)$$

$$\forall y_1 \geq 0, y_2 \leq 0, y_3 \in \mathbb{R} \quad P(y_1, y_2, y_3) \geq \text{opt}(P)$$

(this proves the weak duality thm. by construction)

$\min_{y_1, y_2, y_3} P(y_1, y_2, y_3)$ this will give us the best UB.

$$\begin{aligned} \max \quad & (1 + y_1 + 4y_2 + 2y_3)x_1 + \\ & (-2 + 2y_1 + 3y_2 - y_3)x_2 \\ & (-y_1 + 4y_2 + 2y_3)x_3 \\ & -3y_2 - y_3 \end{aligned}$$

This problem can be solved by inspection.

It is of the form:

$$\max c_1 x_1 + c_2 x_2 + \dots + c_n x_n \quad \left| \begin{aligned} \Rightarrow & \text{if } c_i > 0 \wedge x_i \geq 0 \Rightarrow x_i = +\infty \\ & \text{if } c_i < 0 \wedge x_i \geq 0 \Rightarrow x_i = 0, \text{ bound} \\ & \text{if } c_i \neq 0 \wedge x_i \in \mathbb{R} \Rightarrow \text{unbounded} \\ & \text{if } c_i = 0 \wedge x_i \in \mathbb{R} \Rightarrow \text{bounded} \end{aligned} \right.$$

We are only interested in the cases where the problem is bounded hence:

$$\begin{aligned} \min \quad & -3y_2 - y_3 \\ & y_1 + 4y_2 + 2y_3 = -1 \\ & 2y_1 + 3y_2 - y_3 \leq 2 \\ & -y_1 + 4y_2 + 2y_3 \leq 0 \\ & y_1 \geq 0 \\ & y_2 \leq 0 \\ & y_3 \in \mathbb{R} \end{aligned}$$

$$\begin{aligned} \min \quad & +3y_2 + y_3 \\ & -y_1 + 4y_2 - 2y_3 = -1 \\ & -2y_1 - 3y_2 + y_3 \leq 2 \\ & +y_1 - 4y_2 + 2y_3 \leq 0 \\ & y_1 \leq 0 \\ & y_2 \geq 0 \\ & y_3 \in \mathbb{R} \end{aligned}$$

TASK 6

a) it is a transportation problem;

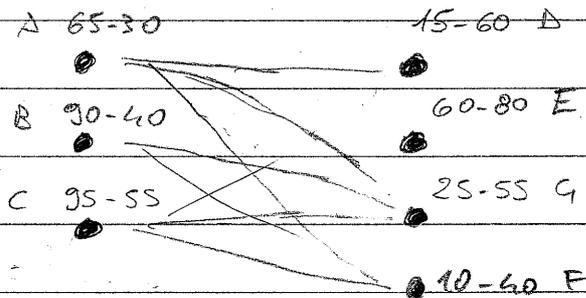
$$\min \sum_{ij} c_{ij} x_{ij}$$

$$\sum_j x_{ij} - \sum_j x_{ji} = f_i - d_i \quad \forall i$$

x_{ij} integer

desired
current

This is a network flow problem,
specifically a transportation problem
hence any alg. for it would go:



↑
those
who
have an
excess of
cars

↑
those
who
require
cars to
be added